Our understanding of life will derive in large measure from *how* we understand chemistry. Walter Fontana and Leo W. Buss, 1996



# Chemistry and Processes

## Luca Cardelli

#### Microsoft Research

The Microsoft Research - University of Trento Centre for Computational and Systems Biology

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www.luca.demon.co.uk/ArtificialBiochemistry.htm

# From Chemistry to ODEs

#### **Chemical Reactions**

| A           | $\rightarrow^{r}$ | $B_1$          | + | ••• | + | B <sub>n</sub>        |
|-------------|-------------------|----------------|---|-----|---|-----------------------|
| $A_1 + A_2$ | $\rightarrow^{r}$ | $B_1$          | + |     | + | <b>B</b> <sub>n</sub> |
| A + A       | $\rightarrow^{r}$ | B <sub>1</sub> | + |     | + | Bn                    |

Degradation Asymmetric Collision Symmetric Collision

[A]• = -r[A]

 $[A_i]^{\bullet} = -r[A_1][A_2]$ 

 $[A]^{\bullet} = -r[A]([A]-1) \qquad Ma$ (assuming  $A \neq B_i \neq A_j$  for all i,j)

Exponential Decay Mass Action Law

Mass Action Law

No other reactions!

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The chemical Langevin equation Daniel T. Gillespie<sup>a)</sup>

Research Department, Code 4T4100D, Naval Air Warfare Center, China Lake, California 93555

Genuinely *trimolecular* reactions do not physically occur in dilute fluids with any appreciable frequency. *Apparently* trimolecular reactions in a fluid are usually the combined result of two bimolecular reactions and one monomolecular reaction, and involve an additional short-lived species.

#### Chapter IV: Chemical Kinetics

[David A. Reckhow, CEE 572 Course] ... reactions may be either elementary or nonelementary. <u>Elementary reactions</u> are those reactions that occur exactly as they are written, without any intermediate steps. These reactions almost always involve just one or two reactants. ... <u>Non-elementary reactions</u> involve a series of two or more elementary reactions. Many complex environmental reactions are non-elementary. In general, reactions with an overall reaction order greater than two, or reactions with some non-integer reaction order are non-elementary.

#### THE COLLISION THEORY OF REACTION RATES

www.chemguide.co.uk

The chances of all this happening if your reaction needed a collision involving more than 2 particles are remote. All three (or more) particles would have to arrive at exactly the same point in space at the same time, with everything lined up exactly right, and having enough energy to react. That's not likely to happen very often!

Trimolecular reactions:  $A + B + C \rightarrow^{r} D$ the measured "r" is an (imperfect) aggregate of e.g.:

 $A + B r^2 \leftrightarrow^{r1} AB$ 

 $AB + C \rightarrow^{r3} D$ 

Enzymatic reactions:

 $S \xrightarrow{E} P$ 

the "r" is given by Michaelis-Menten (approximated steady-state) laws:  $E + S {}^{r2} \leftrightarrow {}^{r1} ES$  $ES \rightarrow {}^{r3} P + E$ 

#### From Reactions to ODE's



#### **Rate Law of Symmetric Reactions**

The proper rate law  $I_4$  of the symmetric reaction is  $k_4[F]([F]-1)/2$  because that is the number of possible collisions between [F] particles.

But the  $v_4$  contribution to [F] is  $-k_4$ [F]([F]-1) because 2 F are consumed in that reaction!

Compare with the contribution of  $v_4$  to [B], which is  $k_4$ [F]([F]-1)/2.

$$[A]^{\bullet} = -I_1 - I_2$$
  

$$[B]^{\bullet} = -I_1 + I_4$$
  

$$[C]^{\bullet} = 2I_1 - I_2 - I_3$$
  

$$[D]^{\bullet} = I_2$$
  

$$[E]^{\bullet} = I_3$$
  

$$[F]^{\bullet} = I_3 - 2I_4$$

$$[A]^{\bullet} = -k_1[A][B] - k_2[A][C]$$
  

$$[B]^{\bullet} = -k_1[A][B] + k_4[F]([F]-1)/2$$
  

$$[C]^{\bullet} = 2k_1[A][B] - k_2[A][C] - k_3[C]$$
  

$$[D]^{\bullet} = k_2[A][C]$$
  

$$[E]^{\bullet} = k_3[C]$$
  

$$[F]^{\bullet} = k_3[C] - k_4[F]([F]-1)$$

That's why we said earlier that  

$$A + A \rightarrow^{r} B_{1} + ... + B_{n}$$
 gives  $(A]^{\bullet} = -r[A]([A]-1)$  (2 A consumed)  
Instead, e.g.:  
 $A + A \rightarrow^{r} A + B$  gives  $[A]^{\bullet} = -r[A]([A]-1)/2$  (1 A consumed)

# From Chemistry to Processes

#### Chemical Ground Form (CGF)

 $\begin{array}{l} \mathsf{E} :::= X_{1} = \mathsf{M}_{1}, \ ..., \ X_{n} = \mathsf{M}_{n} \\ \mathsf{M} :::= \pi_{1} ; \mathsf{P}_{1} \oplus \ ... \oplus \ \pi_{n} ; \mathsf{P}_{n} \\ \mathsf{P} :::= X_{1} \mid ... \mid X_{n} \\ \pi ::= \tau_{r} \ \mathsf{Pn}_{(r)} \ !\mathsf{n}_{(r)} \\ \mathcal{C}GF :::= \mathsf{E}, \mathsf{P} \end{array}$ 

Definitions(n≥0)Molecules(n≥0)Solutions(n≥0)Interactions(delay, input, output)Definitions with Initial Conditions

(To translate chemistry back to processes we need a bit more than simple automata: we may have "+" on the right of  $\rightarrow$ , that is we may need "|" after  $\pi$ .)

 $\oplus$  is stochastic choice (vs. + for chemical reactions) O is the null solution (P|O = O|P = P) and null molecule (M $\oplus$ O = O $\oplus$ M = M) ( $\tau_0$ ;P = O) X<sub>i</sub> are distinct in E Each name n is assigned a fixed rate r: n<sub>(r)</sub>





2 degradations at rate r

[A]• = -r[A] = -2r

(constant half-life!)

 $A = \tau_r; A'$  $A = \tau_r; A'$ 

2 delays at rate r (sampled from exponential distribution)



4 (productive) collisions between 2+2 particles, at rate r 4 interactions between 2+2 processes at rate r

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#### Symmetric Reactions: SPiM vs Matlab

#### $[x1]^{\bullet} = -r[x1]([x1]-1) \iff v@r/2$



#### **Chemical Reactions to CGF**



→ E

## That CGF in SPiM

 $\mathbf{k}_2$ 

 $\mathbf{k}_1$ 



directive sample 10.0
directive plot A(); B(); C(); D(); E(); F()

val k1 = 0.001 new v1@k1:chan val k2 = 0.001 new v2@k1:chan

val k3 = 1.0

val k4 = 0.001 new v4@k4/2.0:chan

let A() = do ?v1;(C()|C()) or ?v2;D()
and B() = !v1
and C() = do !v2 or delay@k3;(E()|F())
and D() = ()

and E() = ()

and F() = do ?v4;B() or !v4

run 300 of (A()|B()|C()|D()|E()|F())

#### ... in other words

From chemical reactions C to a CGF Pi(C):

$$\begin{array}{lll} \mathsf{Pi}(\mathcal{C}) &= \{ (\mathsf{X} = \oplus((\mathsf{v}_i \colon \mathsf{X} \to^{\mathsf{k}} \mathsf{P}) \in \mathcal{C}) \ of(\tau_{\mathsf{k}}; \mathsf{P}) & \oplus \\ & \oplus((\mathsf{v}_i \colon \mathsf{X} + \mathsf{Y} \to^{\mathsf{k}} \mathsf{P}) \in \mathcal{C} \ \text{and} \ \mathsf{Y} \neq \mathsf{X}) \ of(\mathsf{?v}_{\mathsf{i}(\mathsf{k}i)}; \mathsf{P}) & \oplus \\ & \oplus((\mathsf{v}_i \colon \mathsf{Y} + \mathsf{X} \to^{\mathsf{k}} \mathsf{P}) \in \mathcal{C} \ \text{and} \ \mathsf{Y} \neq \mathsf{X}) \ of(\mathsf{!v}_{\mathsf{i}(\mathsf{k}i)}; \mathsf{O}) & \oplus \\ & \oplus((\mathsf{v}_i \colon \mathsf{X} + \mathsf{X} \to^{\mathsf{k}} \mathsf{P}) \in \mathcal{C}) \ of(\mathsf{?v}_{\mathsf{i}(\mathsf{k}i/2)}; \mathsf{P} \oplus \mathsf{!v}_{\mathsf{i}(\mathsf{k}i/2)}; \mathsf{O}) & ) \\ & f(\mathsf{v}_i \colon \mathsf{X} + \mathsf{X} \to^{\mathsf{k}} \mathsf{P}) \in \mathcal{C}) \ of(\mathsf{?v}_{\mathsf{i}(\mathsf{k}i/2)}; \mathsf{P} \oplus \mathsf{!v}_{\mathsf{i}(\mathsf{k}i/2)}; \mathsf{O}) & ) \\ & s.t. \mathsf{X} \ \text{is a species in } \mathcal{C} \} \end{array}$$

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# From Processes Directly to ODEs (hard)

#### "Micromodels": Continuous Time Markov Chains

- The underlying semantics of stochastic  $\pi$ -calculus (and stochastic interacting automata). Well established in many ways.
  - Automata with rates on transitions.
- "The" correct semantics for chemistry, executable.
  - Gillespie stochastic simulation algorithm
- Good description of "individual" behavior
- But does not give a good sense of "collective" properties.
  - Yes one can do simulation.
  - Yes one can do program analysis.
  - Yes one can do modelchecking.
  - But somewhat lacking in "predictive power".

#### "Macromodels": Ordinary Differential Equations

- Micromodels have lots of advantages
  - Compositional, compact, mechanistic, etc.
- But they always ask:
  - "Yes, but how does you automata model relate to the 75 ODE models in the literature?"
- From processes/automata to ODEs directly:
  - In principle: just write down the Rate Equation:
    - Determine the set of all possible *states* S of each process.
    - Determine the rates of the transitions between such states.
    - Let [S] be the "number of processes in state S" as a function of time.
    - Define for each state S:
      - [S] = (rate of change of the number of processes in state S) Cumulative rate of transitions from any state S' to state S, times [S'], minus cumulative rate of transitions from S to any state S", times [S].
  - Intuitive (rate = inflow minus outflow), but often clumsy to write down precisely.
- But why go directly from processes to ODE?
  - If we first convert processes to chemical reactions, then we can convert to ODEs by standard means!

The speed of interaction is proportional to the number of possible interactions.

The rate of change of a quantity is the cumulative result of the (always changing) speed of the interactions that produce it and consume it. It is a **global** and **dynamic** property.



# From Processes to Chemistry

### Automata to Chemistry (by hand)

#### Automata to Chemistry



### (Lack of) Compositionality



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#### Single Automata Models are n<sup>2</sup> More Compact



### **Process Normal Forms to Chemistry**

#### Chemical Parametric Form (CPF)

E ::=  $X_1(\mathbf{p}_1)=M_1, \dots, X_n(\mathbf{p}_n)=M_n$ M ::=  $\pi_1; P_1 \oplus \dots \oplus \pi_n; P_n$ P ::=  $X_1(\mathbf{p}_1) \mid \dots \mid X_n(\mathbf{p}_n)$   $\pi$  ::=  $\tau_r$  ?n(**p**) !n(**p**) CPF::= E,P

CPF is NF without name generation.

CPF is CGF plus parameters.

#### Random example

$$P(a,b) = ?a(c); (P(c,b) | Q(c))$$
  
Q(a) = !a(a); 0  
$$P(x_{(r)},y_{(s)}) | Q(x_{(r)})$$

| Definitions               | (n ≥ 0)        |  |  |
|---------------------------|----------------|--|--|
| Molecules                 | (n ≥ 0)        |  |  |
| Solutions                 | (n ≥ 0)        |  |  |
| Interactions              |                |  |  |
| with initial conditions   |                |  |  |
| stic choice (vs. + for ch | emical poactio |  |  |

 $\oplus$  is stochastic choice (vs. + for chemical reactions) 0 is the null solution (P|0 = 0|P = P) and null molecule (M $\oplus$ 0 = 0 $\oplus$ M = M) ( $\tau_0$ ;P = 0) X<sub>i</sub> are distinct in E, **p** are vectors of names **p** are vectors of distinct names when in binding position Each free name n in E is assigned a fixed rate r: written either n<sub>(r)</sub>, or  $\rho_{CPF}(n)$ =r.

#### Repressilator

 $\begin{array}{l} \left( \mathsf{Neg}(a,b) = ?a; \mathsf{Inh}(a,b) \oplus \tau_{\epsilon}; (\mathsf{Tr}(b) \mid \mathsf{Neg}(a,b)) \right) \\ \mathsf{Inh}(a,b) = \tau_{\eta}; \mathsf{Neg}(a,b) \\ \mathsf{Tr}(b) = !b; \mathsf{Tr}(b) \oplus \tau_{\delta}; \mathsf{O} \\ \mathsf{Neg}(x,y) \mid \mathsf{Neg}(y,z) \mid \mathsf{Neg}(z,x) \end{array} \right)$ 

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#### CPF to Chemistry: the CGF Case

CGF = The CPF subset with no communication, no parametric definitions.

| $E ::= X_1 = M_1,, X_n = M_n$               | Definitions                         | (n≥0)                    |
|---|-------------------------------------|--------------------------|
| $M ::= \pi_1; P_1 \oplus \oplus \pi_n; P_n$ | Molecules                           | (n≥0)                    |
| $P ::= X_1 \mid \dots \mid X_n$             | Solutions                           | (n≥0)                    |
| $\pi ::= \tau_r ?n_{(r)} !n_{(r)}$          | Interactions                        | s (delay, input, output) |
| <i>CG</i> F ::= E,P                         | Definitions with Initial Conditions |                          |

Each X in E is seen as a separate *species*.

Chemical reactions for E:  $(N.B.: \{...\}^m$  is a multiset, and P is P with all the | changed to +)

$$\begin{array}{ll} Ch_{\mathcal{G}}(E) \coloneqq \{(X \rightarrow^{r} P) \ s.t. \ (X \equiv \tau_{r}; P \oplus ...) \in E\}^{m} \\ \cup^{m} \{(X + Y \rightarrow^{r} P + Q) \ s.t. \ X \neq Y, \ \langle (X \equiv 2n_{(r)}; P \oplus ...), (Y \equiv !n_{(r)}; Q \oplus ...) \rangle \in E^{2}\}^{m} \\ \cup^{m} \{(X + X \rightarrow^{2r} P + Q) \ s.t. \ (X \equiv 2n_{(r)}; P \oplus ... \equiv !n_{(r)}; Q \oplus ...) \rangle \in E\}^{m} \end{array}$$

Initial conditions for P:



### (Note on computing the multisets)

A multiset  $M \in Multiset(S)$ , where S is a set with equality, is a total function  $S \rightarrow Nat$ , which may also be written as a finite enumeration with repetitions:  $\{...\}^m$ .

Multiset binary union is the function  $\cup^{m}(M,M') = \s. M(s)+M'(s)$ . Multiset big union over a finite index set S is written:  $\cup^{m} i \in S \ of M_{i}$ .

- The shorthand: {(X  $\rightarrow$ <sup>r</sup> P) s.t. (X =  $\tau_r$ ; P  $\oplus$  ...)  $\in$  E}<sup>m</sup>

means: "for each  $(X=\pi_1; P_1 \oplus ... \oplus \pi_n; P_n) \in E$  and for each i such that  $\pi_i = \tau_r$ , return a copy of  $(X \rightarrow P_i)$ ".

i.e., it is defined as the following finite union of singleton multisets:

 $\cup^{\mathsf{m}}\{(\mathsf{X}=\pi_1;\mathsf{P}_1\oplus...\oplus\pi_n;\mathsf{P}_n)\in\mathsf{E}\}\quad \text{ of }(\cup^{\mathsf{m}}\{\mathsf{i} \ \mathsf{s}.\mathsf{t}.\ \pi_\mathsf{i}=\tau_\mathsf{r}\}\quad \text{ of }\{(\mathsf{X}\to^\mathsf{r}\mathsf{P}_\mathsf{i})\}^{\mathsf{m}})$ 

- The shorthand:  $\{(X + Y \rightarrow^r P + Q) \text{ s.t. } X \neq Y, \langle (X \equiv ?n_{(r)}; P \oplus ...), (Y \equiv !n_{(r)}; Q \oplus ...) \rangle \in E^2\}^m$ 

means: "for each ordered pair  $(X=\pi_1; P_1 \oplus ... \oplus \pi_n; P_n), (Y=\sigma_1; Q_1 \oplus ... \oplus \sigma_n; Q_n)$  from  $E^2$  with  $X \neq Y$ , for each  $n_{(r)}$ , for each i such that  $\pi_i = 2n_{(r)}$ , and for each j such that  $\sigma_i = 2n_{(r)}$ , return a copy of the reaction  $(X + Y \rightarrow r P_i + Q_i)$ "

i.e., it is defined as the following finite union of singleton multisets:

 $\bigcup_{m \in \{X=\pi_1; P_1 \oplus ... \oplus \pi_n; P_n\}, (Y=\rho_1; Q_1 \oplus ... \oplus \rho_m; Q_m) \in E^2 \text{ with } X \neq Y }$ of  $(\bigcup_{i,j>s,t, \pi_i=2n_{(r)}, \rho_j=!n_{(r)}}$ of  $\{(X + Y \rightarrow^r P_i + Q_j)\}^m$ 

- The shorthand:  $\{(X + X \rightarrow^{2r} P + Q) \text{ s.t. } (X \equiv ?n_{(r)}; P \oplus ... \equiv !n_{(r)}; Q \oplus ...) \in E\}^m$ 

means: "for each  $(X=\pi_1; P_1 \oplus ... \oplus \pi_n; P_n)$  in E, for each  $n_{(r)}$ , for each i such that  $\pi_i = 2n_{(r)}$ , and for each j such that  $\pi_j = 1n_{(r)}$ , return a copy of the reaction  $(X + X \rightarrow 2^r P_i + P_j)$ "

#### Example





## Subtler Examples

These are not *finite state systems*, but *finite species systems* are ok!

Unbounded state, but only 1 species. No problem!

#### Multisets:

The same interaction can occur multiple times and must be taken into account:

E:  

$$A = 2n; B \oplus 2n; B$$
  
 $C = 1n; D$   
 $C = 2n; B \oplus 2n; B$   
 $A + C \rightarrow^{p(n)} B + D$   
 $A + C \rightarrow^{p(n)} B + D$ 

$$\begin{array}{c} \text{Fhat is:} \\ A + C \rightarrow^{2\rho(n)} B + D \end{array}$$

Symmetric reactions:



The rate of a was pre-halved and must be restored.

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## CPF to Chemistry: Handling Parameters

Consider first the CPF subset with no communication (pure ?n, !n).

**Grounding** (replace parameters with constants) where X/p is a name in bijection with  $\langle X, p \rangle$ (each X/p is seen as a separate *species*)  $/(\pi_1; P_1 \oplus ... \oplus \pi_n; P_n) =_{def} \pi_1; /(P_1) \oplus ... \oplus \pi_n; /(P_n)$  $/(X_1(p_1) | ... | X_n(p_n)) =_{def} X_1/p_1 | ... | X_n/p_n$   $E ::= X_{1}(\mathbf{p}_{1}) = M_{1}, ..., X_{n}(\mathbf{p}_{n}) = M_{n}$  $M ::= \pi_{1}; P_{1} \oplus ... \oplus \pi_{n}; P_{n}$  $P ::= X_{1}(\mathbf{p}_{1}) | ... | X_{n}(\mathbf{p}_{n})$  $\pi ::= \tau_{r} ?n !n$ 

Let N be the set of free names occurring in E (names not bound by definitions).

 $E_G$  is the **Parametric Explosion** of E (still a finite species system) computed by replacing parameters with all combinations of free names in E

 $E_{G} := \{(X/q = /(M\{p \leftarrow q\})) \text{ s.t. } (X(p) = M) \in E \text{ and } q \in N^{\#p}\}$   $P_{G} := /P \qquad (\text{simply ground the given initial conditions once})$   $E_{G} \text{ is a } CGF! \text{ To obtain the chemical reactions } Ch_{P}(E), \text{ just compute } Ch_{G}(E_{G})$   $Ch_{P}(E) = Ch_{G}(E_{G})$ 



Note that  $E_G$  is not necessarily the "full" explosion (there may not be enough free variables in E,P to generate it). In particular, if E,P is closed (no free variables), then  $E_G$  contains only the parameterless definitions, which must then be closed and hence unrelated to any parametric definitions. One can add initial conditions  $X_1(\mathbf{p}_1) \mid ... \mid X_n(\mathbf{p}_n)$  where  $X_i$  are all the definitions in E and  $\mathbf{p}_1 ... \mathbf{p}_n$  are made of distinct (free) variables. Then  $E_G$  becomes the most general set of chemical reactions from E, out of any possible initial conditions.

## **Iterative Algorithm**

The chemical reactions for E, computed directly from the parametric explosion  $E_G$  are highly redundant because they include all the parameter permutation symmetries. The following iterative algorithm, for the parametric case, computes a subset of  $E_G$  from the initial conditions of E,P. It produces a (usually) much smaller although not necessarily minimal set C.

Here:

- C is the incrementally built set of chemical reactions
- $E_c$  (a CGF!) is the incrementally built set of grounded definitions from the original E.

 $\begin{array}{l} \mbox{initialization} \\ E_{\mathcal{C}} := \{ (X/q = /(M\{p \leftarrow q\})) \ \mbox{s.t. } X(q) \ \mbox{occurs in } P \ \mbox{and } (X(p) = M) \in E \} \ \ \mbox{(from initial conditions } P) \\ \mbox{iteration} \\ \mathcal{C} := & Ch_{\mathcal{G}}(E_{\mathcal{C}}) \\ E_{\mathcal{C}}' := & E_{\mathcal{C}} \cup \{ (X/q = /(M\{p \leftarrow q\})) \ \mbox{s.t. } X/q \ \mbox{occurs in } C \ \mbox{and } (X(p) = M) \in E \} \\ \mbox{termination} \\ \mbox{if } E_{\mathcal{C}}' = E_{\mathcal{C}} \ \mbox{then stop and return } (\mathcal{C},/(P)) \ \mbox{else } E_{\mathcal{C}} := E_{\mathcal{C}}' \ \mbox{and iterate.} \end{array}$ 

The algorithm terminates because  $E_c$  never shrinks and is always a subset of  $E_c$ , which is finite.

# Example: Neg(x,x)

E =

Neg(a,b) = ?a; Inh(a,b) ⊕ τ₂; (Tr(b) | Neg(a,b)) Inh(a,b) =  $\tau_n$ ; Neg(a,b) Tr(b) = !b; Tr(b)  $\oplus \tau_s$ ; 0 Neq(x,x)---- initialization ----- $E_{c} := \{ Neg/x, x = ?x; Inh/x, x \oplus \tau_{c}; (Tr/x | Neg/x, x) \}$ ----- iteration 1 ----- $C := \{ Neg/x, x \rightarrow^{\varepsilon} Tr/x + Neg/x, x \}$  $E_{x} = \{ Neg/x, x = ?x; Inh/x, x \oplus \tau_{x}; (Tr/x | Neg/x, x) \}$ Tr/x = !x; Tr/x  $\oplus \tau_s$ ; 0} ----- iteration 2 ----- $C := \{ Neg/x, x \rightarrow^{\varepsilon} Tr/x + Neg/x, x \}$  $Tr/x \rightarrow \delta 0$  $Tr/x + Neg/x, x \rightarrow \rho(x) Tr/x + Inh/x, x$  $E_{x} = \{ Neg/x, x = ?x; Inh/x, x \oplus \tau_{x}; (Tr/x | Neg/x, x) \}$ Tr/x = !x; Tr/x  $\oplus \tau_s$ ; 0 Inh/x,x =  $\tau_n$ ; Neg/x,x}

----- iteration 3 -----  $C := \{ Neg/x, x \rightarrow^{\varepsilon} Tr/x + Neg/x, x$  $Tr/x \rightarrow^{\delta} 0$  $Tr/x + Neg/x, x \rightarrow^{\rho(x)} Tr/x + Inh/x, x$  $Inh/x, x \rightarrow^{\eta} Neg/x, x \}$  $E_{c} := no change$ ----- termination ----- $Neg/x, x \rightarrow^{\varepsilon} Tr/x + Neg/x, x$  $Tr/x \rightarrow^{\delta} 0$  $Tr/x + Neg/x, x \rightarrow^{\rho(x)} Tr/x + Inh/x, x$  $Inh/x, x \rightarrow^{\eta} Neg/x, x$ Neg/x, xNeg/x, xNeg/x

# Example: Tr(x) | Neg(x,y)

E =

----- iteration 2 ----- $C := \{ \operatorname{Tr}/x \to^{\delta} 0 \\ \operatorname{Neg}/x, y \to^{\varepsilon} \operatorname{Tr}/y + \operatorname{Neg}/x, y \\ \operatorname{Tr}/x + \operatorname{Neg}/x, y \to^{\rho(x)} \operatorname{Tr}/x + \operatorname{Inh}/x, y \\ \operatorname{Tr}/y \to^{\delta} 0 \\ \operatorname{Inh}/x, y \to^{\eta} \operatorname{Neg}/x, y \}$   $E^{c}:= \text{ no change}$ ----- termination ----- $\operatorname{Tr}/x \to^{\delta} 0 \\ \operatorname{Neg}/x, y \to^{\varepsilon} \operatorname{Tr}/y + \operatorname{Neg}/x, y \\ \operatorname{Tr}/x + \operatorname{Neg}/x, y \to^{\rho(x)} \operatorname{Tr}/x + \operatorname{Inh}/x, y \\ \operatorname{Tr}/y \to^{\delta} 0 \\ \operatorname{Inh}/x, y \to^{\eta} \operatorname{Neg}/x, y \\ \operatorname{Tr}/x + \operatorname{Neg}/x, y$ 

## **CPF to Chemistry: Handling Communication**

**Grounding** (replace parameters with constants) just one main change: now also convert each input parameter into a ground choice of all possible inputs N is the set of free names in E.P. **#p** is the length of **p** n/p is a name in bijection with <n,p> X/p is a name in bijection with  $\langle X, p \rangle$ (each X/p is seen as a separate *species*)  $/_{N}(\tau_{r};P) = \tau_{r}; /_{N}(P)$  $/_{N}(!n_{(r)}(\mathbf{p});P) = !n/\mathbf{p}_{(r)}; /_{N}(P)$  $/_{N}(?n_{(r)}(\mathbf{p});P) = \bigoplus (\mathbf{q} \in N^{\#p}) \text{ of } ?n/\mathbf{q}_{(r)}; /_{N}(P\{\mathbf{p} \leftarrow \mathbf{q}\})$  $/_{N}(\pi_{1};\mathsf{P}_{1}\oplus...\oplus\pi_{n};\mathsf{P}_{n})=/_{N}(\pi_{1};\mathsf{P}_{1})\oplus...\oplus/_{N}(\pi_{n};\mathsf{P}_{n})$  $/_{N}(X_{1}(\mathbf{p}_{1}) \mid ... \mid X_{n}(\mathbf{p}_{n})) = X_{1}/\mathbf{p}_{1} \mid ... \mid X_{n}/\mathbf{p}_{n}$  $E_{G}$  is again the **Parametric Explosion** of E  $E_G := \{ (X/q = /_N(M\{p \leftarrow q\})) \text{ s.t. } (X(p) = M) \in E \text{ and } q \in N^{\#p} \}$  $P_G := /_N(P)$ (simply ground the given initial conditions once)

 $Ch(E) = Ch_{G}(E_{G})$   $E_{G}$  is a again a CGF!

 $E ::= X_{1}(\mathbf{p}_{1}) = M_{1}, ..., X_{n}(\mathbf{p}_{n}) = M_{n}$   $M ::= \pi_{1}; P_{1} \oplus ... \oplus \pi_{n}; P_{n}$   $P ::= X_{1}(\mathbf{p}_{1}) | ... | X_{n}(\mathbf{p}_{n})$  $\pi ::= \tau_{r} ?n(\mathbf{p}) !n(\mathbf{p})$ 

#### Example E (with free names x,y): $Q/x + P/x, y \rightarrow^{\rho(x)} P/x, y + Q/x$ P(a,b) = 2a(c); (P(c,b) | Q(c)) $Q/x + P/x, x \rightarrow^{\rho(x)} P/x, x + Q/x$ Q(a) = !a(a); 0 $Q/y + P/y, x \rightarrow^{\rho(y)} P/y, x + Q/y$ $P(x,y) \mid Q(x)$ $Q/y + P/y, y \rightarrow^{\rho(y)} P/y, y + Q/y$ P/x,y + Q/xParametric Explosion $E_G$ : Reduced C **Chemical Reactions** Q/x = !x/x; 0 (from initial conditions) Q/y = !y/y; 0P/x,y = **?**x/x; (P/x,y | Q/x) ⊕ $Q/x + P/x, y \rightarrow^{p(x)} P/x, y + Q/x$ **?x/y**; (P/y,y | Q/y) P/x,y + Q/xP/y,x = ?y/x; (P/x,x | Q/x) ⊕ ?y/y; (P/y,x | Q/y) P/x,x = ?x/x; (P/x,x | Q/x) ⊕ ?x/y; (P/y,x | Q/y) P/y,y = ?y/x; (P/x,y | Q/x) ⊕ ?y/y; (P/y,y | Q/y) P/x,y | Q/x

### Example

E (with free names x,y): P(a,b) = ?a(c); (P(c,b) | Q(c,c))  $Q(a,b) = !a(a); 0 \oplus !b(b); 0$ P(x,y) | Q(x,y)

 $E_G$ : Q/x,y= !x/x; 0 ⊕ !y/y; 0  $Q/y,x = \frac{1}{y}/y; 0 \oplus \frac{1}{x}/x; 0$ Q/x,x= !x/x; 0 ⊕ !x/x; 0 Q/y,y = !y/y; 0 ⊕ !y/y; 0 P/x,y = **?**x/x; (P/x,y | Q/x,x) ⊕ ?x/y; (P/y,y | Q/y,y) P/y,x = ?y/x; (P/x,x | Q/x,x) ⊕ ?y/y; (P/y,x | Q/y,y) P/x,x = ?x/x; (P/x,x | Q/x,x) ⊕ ?x/y; (P/y,x | Q/y,y) P/y,y = ?y/x; (P/x,y | Q/x,x)⊕ ?y/y; (P/y,y | Q/y,y) P/x,y | Q/x

*C*:  $Q/x,y + P/x,y \rightarrow^{\rho(x)} P/x,y + Q/x,x$  $Q/x,y + P/x,x \rightarrow \rho(x) P/x,x + Q/x,x$  $Q/x,y + P/y,x \rightarrow^{\rho(y)} P/y,x + Q/y,y$  $Q/x,y + P/y,y \rightarrow^{\rho(y)} P/y,y + Q/y,y$  $Q/y,x + P/y,x \rightarrow \rho(y) P/y,x + Q/y,y$  $Q/y, x + P/y, y \rightarrow^{\rho(y)} P/y, y + Q/y, y$  $Q/y,x + P/x,y \rightarrow^{\rho(x)} P/x,y + Q/x,x$  $Q/y,x + P/x,x \rightarrow \rho(x) P/x,x + Q/x,x$  $Q/x,x + P/x,y \rightarrow^{2\rho(x)} P/x,y + Q/x,x$  $Q/x,x + P/x,x \rightarrow 2\rho(x) P/x,x + Q/x,x$  $Q/y,y + P/y,x \rightarrow^{2\rho(y)} P/y,x + Q/y,y$  $Q/y,y + P/y,y \rightarrow 2\rho(y) P/y,y + Q/y,y$ P/x,y + Q/x,y

### **Exponential Explosion via Communication**

Parametric form (size 3).

$$X(t,f) = !t(f); X(f,t)$$

$$A(x_1,x_2,x_3) =$$

$$?x_1(y); A(y,x_2,x_3) \oplus$$

$$?x_2(y); A(x_1,y,x_3) \oplus$$

$$?x_3(y); A(x_1,x_2,y)$$

$$A(t_1,t_2,t_3) |$$

$$X(t_1,f_1) | X(t_2,f_2) | X(t_3,f_3)$$

green = free names with associated rates

Exponentially larger ground form (size 8).

X/t<sub>1</sub>,f<sub>1</sub> = !t<sub>1</sub>/f<sub>1</sub>; X/f<sub>1</sub>,t<sub>1</sub>  
X/f<sub>1</sub>,t<sub>1</sub> = !f<sub>1</sub>/t<sub>1</sub>; X/f<sub>1</sub>,t<sub>1</sub>  
X/t<sub>2</sub>,f<sub>2</sub> = !t<sub>2</sub>/f<sub>2</sub>; X/f<sub>2</sub>,t<sub>2</sub>  
X/f<sub>2</sub>,t<sub>2</sub> = !f<sub>2</sub>/t<sub>2</sub>; X/t<sub>2</sub>,f<sub>2</sub>  
X/t<sub>3</sub>,f<sub>3</sub> = !t<sub>3</sub>/f<sub>3</sub>; X/f<sub>3</sub>,t<sub>3</sub>  
X/f<sub>3</sub>,t<sub>3</sub> = !t<sub>3</sub>/f<sub>3</sub>; X/f<sub>3</sub>,t<sub>3</sub>  
A/t<sub>1</sub>,t<sub>2</sub>,t<sub>3</sub> =  
$$?t_1/f_1; A/f_1, t_2, t_3 \oplus$$
  
 $?t_2/f_2; A/t_1, f_2, t_3 \oplus$   
 $?t_3/f_3; A/t_1, t_2, f_3$   
A/f<sub>1</sub>,t<sub>2</sub>,t<sub>3</sub> =  
 $?f_1/t_1; A/t_1, t_2, t_3 \oplus$   
 $?t_2/f_2; A/t_1, f_2, t_3 \oplus$   
 $?t_2/f_2; A/t_1, f_2, t_3 \oplus$   
 $?t_3/f_3; A/t_1, t_2, f_3$   
...etc for the other bit flips  
A/t<sub>1</sub>,t<sub>2</sub>,t<sub>3</sub> |  
X/t<sub>1</sub>,t<sub>1</sub> | X/t<sub>2</sub>,f<sub>2</sub> | X/t<sub>3</sub>,f<sub>3</sub>

We have not fully expanded the input summations in A(...): only enough for the channels that are actually used: e.g., there would never be a communication on channel  $t_1/f_2$ .

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## **CPF Iterative Algorithm**

The same as before, just using the new grounding  $/_{N}$  (based on the set N of free names of E,P, for expanding inputs).

 $\begin{array}{l} \mbox{initialization} \\ N = fn(E,P) \\ E_{c} := \{(X/q = /_{N}(M\{p \leftarrow q\})) \ s.t. \ X(q) \ occurs \ in \ P \ and \ (X(p) = M) \in E\} \ (initial \ conditions) \\ \mbox{iteration} \\ C := Ch_{G}(E_{c}) \\ E_{c}' := E_{c} \cup \{(X/q = /_{N}(M\{p \leftarrow q\})) \ s.t. \ X/q \ occurs \ in \ C \ and \ (X(p) = M) \in E\} \\ \mbox{itermination} \\ \mbox{if } E_{c}' = E_{c} \ then \ stop \ and \ return \ (C, /_{N}(P)) \ else \ E_{c} := E_{c}' \ and \ iterate. \end{array}$ 

#### A General $\pi$ -calculus Normal Form

| E  | $ ::= X_1(\mathbf{p}_1) = (v\mathbf{q}_1)M_1,, X_n(\mathbf{p}_n) = (v\mathbf{q}_n)M_n $ |
|----|---|
| Μ  | $::= \pi_1; P_1 \oplus \oplus \pi_n; P_n$   |
| Ρ  | $::= X_1(\mathbf{p}_1) \mid \mid X_n(\mathbf{p}_n)$                                     |
| π  | $:= \tau_{r} ?n(p) !n(p)$   |
| NF | ::= E,P   |

| Definitions             | (n ≥ 0) |  |
|-------------------------|---------|--|
| Molecules               | (n ≥ 0) |  |
| Solutions               | (n ≥ 0) |  |
| Interactions            |         |  |
| with initial conditions |         |  |

Any  $\pi$ -calculus process can be written in this normal form. But not all of those can be translated to chemical reactions, because we need to restrict to a *finite number of species*.

Example (complexation):

$$P(a) = (vn_{(d)}) !a(n); P'(a,n)$$

$$P'(a,n) = !n; P(a)$$

$$Q(a) = ?a(n); Q'(a,n)$$

$$Q'(a,n) = ?n; Q(a)$$

$$P(x_{(c)}) | Q(x_{(c)})$$

Example (enzymatic reaction):

 $E(a) = (vn_{(d)}, m_{(p)}) !a(n,m); E'(a,n,m)$   $E'(a,n,m) = !n; E(a) \oplus !m; E(a)$  S(a) = ?a(n,m); S'(a,n,m)  $S'(a,n,m) = ?n; S(a) \oplus ?m; P()$   $P() = \tau_0; 0$  $E(x_{(c)}) | S(x_{(c)}) | S(x_{(c)})$ 

#### Summary

- From Chemistry to ODEs
  - Via the stoichiometric matrix
- From Chemistry to Processes
  - Via the interaction matrix
- From Processes to Chemistry
  - Ground processes: by analysis of normal form interactions
  - Parametric processes: by parametric explosion

#### • Compositionality

- Processes are compositional in components (you just "add components") and chemistry is compositional in reactions (you just "add reactions"). But:
  - Adding reactions has a *linear* effect on the number of components (adding one reaction adds one or two more process actions).
  - Adding components has a *quadratic* effect on the number of reactions (adding one species may add one reaction for each species).
  - ODEs grow quadratically along with the chemical reactions (stoichiometric matrix has one column per reaction).
- Parameterization is available only on processes.

