# Chemistry and Processes 

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## From Chemistry to ODEs

## Chemical Reactions

$$
\begin{array}{lllll}
A & \rightarrow^{r} B_{1}+\ldots+B_{n} & \text { Degradation } & {[A]^{0}=-r[A]} & \text { Exponential Decay } \\
A_{1}+A_{2} \rightarrow r B_{1}+\ldots+B_{n} & \text { Asymmetric Collision } & {\left[A_{i}\right]^{0}=-r\left[A_{1}\right]\left[A_{2}\right]} & \text { Mass Action Law } \\
A+A & \rightarrow^{r} B_{1}+\ldots+B_{n} & \text { Symmetric Collision } & {[A]^{\circ}=-r[A]([A]-1)} & \text { Mass Action Law } \\
& & \text { (assuming } A \neq B_{i} \neq A_{j} \text { for all } i, j \text { ) }
\end{array}
$$

## No other reactions!

| Jourval of Chemical physics volume 113, NUMBER 1 | Chapter IV: Chemical Kinetics | THE COLLISION THEORY OF |
| :---: | :---: | :---: |
| The chemical Langevin equation | [David A. Reckhow , CEE 572 Course] <br> reactions may be either elementary or non- | REACTION RATES <br> www.chemguide.co.uk |
| Daniel T. Gillespie ${ }^{\text {a) }}$ <br> Research Department, Code 4T4100D, Naval Air Warfare Center, China Lake, California 93555 | elementary. Elementary reactions are those reactions that occur exactly as they are written, without any intermediate steps. These reactions almost always | The chances of all this happening if your reaction needed a collision involving more than 2 particles are remote. All three (or |
| Genuinely trimolecular reactions do not physically occur in dilute fluids with any appreciable frequency. Apparently | involve just one or two reactants. ... Non-elementary reactions involve a series of two or more elementary | more) particles would have to arrive at exactly the same point in space at the same |
| trimolecular reactions in a fluid are usually the combined result of two bimolecular reactions and one monomolecular reaction, and involve an additional short-lived species. | reactions. Many complex environmental reactions are non-elementary. In general, reactions with an overall reaction order greater than two, or reactions with some non-integer reaction order are non-elementary. | time, with everything lined up exactly right, and having enough energy to react. That's not likely to happen very often! |

## Trimolecular reactions:

$$
A+B+C \rightarrow^{r} D
$$

the measured " $r$ " is an (imperfect) aggregate of e.g.:

$$
A+B^{r 2} \leftrightarrow r^{r 1} A B
$$

$$
A B+C \rightarrow \rightarrow_{0}^{r 3} D
$$

## Enzymatic reactions:

$$
S \xrightarrow{E} P
$$

the " $r$ " is given by Michaelis-Menten: (approximated steady-state) laws:
$E+S^{r 2} \leftrightarrow r^{r 1} E S$
ES. $\rightarrow$ r3 P+. E

## From Reactions to ODE's

$$
\begin{aligned}
& v_{1}: A+B \rightarrow k_{1} C+C \\
& v_{2}: A+C \rightarrow k_{2} D \\
& v_{3}: C \rightarrow k_{3} E+F \\
& v_{4}: F+F \rightarrow k_{4} B \\
& \begin{array}{c}
\text { Quantity } \\
\text { changes }
\end{array} \\
& \begin{array}{c}
\text { Stoichiometric } \\
\text { matrix }
\end{array} \\
& {[X]^{\bullet}=\mathbf{N} \cdot I}
\end{aligned}
$$

$[A]^{0}=-I_{1}-I_{2}$
$[B]^{\circ}=-l_{1}+l_{4}$
$[C]^{\circ}=2 I_{1}-I_{2}-I_{3}$
[D] ${ }^{0}=I_{2}$
$[\mathrm{E}]^{0}=\mathrm{I}_{3}$
$[\mathrm{F}]^{\circ}=\mathrm{I}_{3}-2 \mathrm{I}_{4}$

Write the coefficients by columns

|  | reactions |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { n } \\ & . \frac{\alpha}{u} \\ & \text { ù } \\ & \text { n } \end{aligned}$ | N | $\mathrm{v}_{1}$ | $\mathrm{v}_{2}$ | $\mathrm{v}_{3}$ | $\mathrm{V}_{4}$ |
|  | A | -1 | -1 |  |  |
|  | B | -1 |  |  | 1 |
|  | C | 2 | -1 | -1 |  |
|  | D |  | 1 |  |  |
|  | E |  |  | 1 |  |
|  | F |  |  | 1 | -2 |
|  | X |  |  |  |  |

Read the concentration changes from the rows
E.g. $[\mathbf{A}]^{\circ}=$
$-k_{1}[A][B]-k_{2}[A][C]$

Set a rate law for each reaction (Degradation/Asymmetric/Symmetric)

| 1 |  |
| :---: | :---: |
| $I_{1}$ | $\mathrm{k}_{1}[\mathbf{A}][\mathrm{B}]$ |
| $\mathrm{I}_{2}$ | $\mathrm{k}_{2}[\mathbf{A}][\mathrm{C}]$ |
| $\mathrm{I}_{3}$ | $\mathrm{k}_{3}[\mathbf{C}]$ |
| $\mathrm{I}_{4}$ | $\mathrm{k}_{4}[\mathrm{~F}]([\mathrm{F}]-1) / 2$ |

$X$ : chemical species
[-]: quantity of molecules
I: rate laws
k: kinetic parameters
N : stoichiometric matrix


## Rate Law of Symmetric Reactions

The proper rate law $\mathrm{I}_{4}$ of the symmetric reaction is $k_{4}[F]([F]-1) / 2$ because that is the number of possible collisions between [F] particles.

But the $v_{4}$ contribution to $[F]^{\cdot}$ is
$-k_{4}[F]([F]-1)$ because $2 F$ are consumed in that reaction!

$$
\begin{aligned}
& {[A]^{\bullet}=-I_{1}-I_{2}} \\
& {[B]^{\bullet}=-I_{1}+I_{4}} \\
& {[C]^{\bullet}=2 I_{1}-I_{2}-I_{3}} \\
& {[D]^{\bullet}=I_{2}} \\
& {[E]^{\bullet}=I_{3}} \\
& {[F]^{\circ}=I_{3}-2 I_{4}}
\end{aligned}
$$

$$
[A]^{\cdot}=-k_{1}[A][B]-k_{2}[A][C]
$$

$$
[B]^{\circ}=-k_{1}[A][B]+k_{4}[F]([F]-1) / 2
$$

$$
[C]^{\bullet}=2 \mathrm{k}_{1}[A][B]-\mathrm{k}_{2}[A][C]-\mathrm{k}_{3}[C]
$$

$$
[D]^{\circ}=k_{2}[A][C]
$$

$$
[E]^{\circ}=k_{3}[C]
$$

$$
[F]^{0}=\mathrm{k}_{3}[C]-\mathrm{k}_{4}[\mathrm{~F}]([\mathrm{F}]-1)
$$

That's why we said earlier that
$A+A \rightarrow^{r} B_{1}+\ldots+B_{n}$ gives $[A]^{\circ}=-r[A]([A]-1) \quad(2 A$ consumed $)$
Instead, e.g.:
$A+A \rightarrow \rightarrow^{r} A+B \quad$ gives $[A]^{\circ}=-r[A]([A]-1) / 2$ (1 A consumed)

## From Chemistry to Processes

## Chemical Ground Form (CGF)

$E::=X_{1}=M_{1}, \ldots, X_{n}=M_{n} \quad$ Definitions ( $n \geq 0$ )
$M::=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n} \quad$ Molecules $\quad(n \geq 0)$
$\mathrm{P}::=\mathrm{X}_{1}|\ldots| \mathrm{X}_{\mathrm{n}}$
$\pi::=\tau_{r} ? n_{(r)}!n_{(r)}$
CGF :: = E, P
Solutions ( $n \geq 0$ )
Interactions (delay, input, output)
Definitions with Initial Conditions
(To translate chemistry back to processes we need a bit more than simple automata: we may have "+" on the right of $\rightarrow$, that is we may need "|" after $\pi$.)

```
\(\oplus\) is stochastic choice (vs. + for chemical reactions)
0 is the null solution ( \(P|0=0| P=P\) )
    and null molecule \((M \oplus 0=0 \oplus M=M)\left(\tau_{0} ; P=0\right)\)
\(X_{i}\) are distinct in \(E\)
    Each name \(n\) is assigned a fixed rate \(r: n_{(r)}\)
```



## Coding of Degradation Reactions ( $A \rightarrow{ }^{r} A^{\prime}$ )

Particles
Rate Law
Processes
$[A]=1$


$$
\begin{aligned}
{[A]^{\bullet} } & =-r[A] \\
& =-r
\end{aligned}
$$

1 degradation at rate $r$
$[A]=2$


2 degradations at rate $r$
$[A]^{\bullet}=-r[A]$
$=-2 r$
(constant half-life!)

$$
\begin{aligned}
& A=\tau_{r} ; A^{\prime} \\
& A=\tau_{r} ; A^{\prime}
\end{aligned}
$$

2 delays at rate $r$
(sampled from exponential distribution)
(or rather, the CGF:

$$
A=\tau_{r} ; A^{\prime}
$$

$$
A \mid A
$$

)

## Coding of Asymmetric Reactions ( $A+B \rightarrow{ }^{r} A^{\prime}+B^{\prime}$ )

Particles
$[A]=1$
$[B]=1$


1 (productive) collision between 1+1 particles, at rate $r$

Rate Law
$\begin{aligned} {[A]^{\circ} } & =-r[A][B] \\ & =-r\end{aligned}$

$$
=-r
$$

$$
B=!v_{(r)} ; B^{\prime}
$$

1 interactions between $1+1$ processes at rate $r$
[A]=2
$[B]=2$


4 (productive) collisions
between $2+2$ particles, at rate $r$
$[A]^{\circ}=-r[A][B]$
$=-4 r$

4 interactions between $2+2$ processes at rate $r$

## Coding of Symmetric Reactions ( $A+A \rightarrow A^{r}+A^{\prime \prime}$ )

Particles
Rate Law

$$
\begin{aligned}
{[A]^{\circ} } & =-r[A]([A]-1) \\
& =0
\end{aligned}
$$

$[A]=2$


1 collision between 2 particles at rater
$[A]^{0}=-r[A]([A]-1)$
$=-2 r$
(remember: 2 A are removed per collision assuming $A^{\prime}, A^{\prime \prime} \neq A$ )


$$
A=? v_{(r / 2)} ; A^{\prime} \oplus!v_{(r / 2)} ; A^{\prime \prime}
$$

0 interactions between 1 process

$$
A=? v_{(r / 2)} ; A^{\prime} \oplus!v_{(r / 2)} ; A^{\prime \prime}
$$

2 interactions between 2 processes at rate $\mathrm{r} / 2$

3 collisions at rate $r$

$$
\begin{aligned}
{[A]^{\bullet} } & =-r[A]([A]-1) \\
& =-6 r
\end{aligned}
$$

3 collisions between 3 particles at rate $r$


6 interactions between 3 processes at rate r/2

## Symmetric Reactions: SPiM vs Matlab

$$
[x 1]^{\circ}=-r[x 1]([x 1]-1) \Leftrightarrow v @ r / 2
$$




## Chemical Reactions to CGF

$$
\begin{aligned}
& \mathrm{v}_{1}: A+B \rightarrow \mathrm{k}_{1} C+C \\
& \mathrm{v}_{2}: A+C \rightarrow \mathrm{k}_{2} D \\
& \mathrm{v}_{3}: C \rightarrow \mathrm{k}_{3} E+F \\
& \mathrm{v}_{4}: F+F \rightarrow \mathrm{k}_{4} B
\end{aligned}
$$

Interaction Matrix

|  | channels and rates <br> (1 per reaction) |  |  |  |  | Half-rate for symmetric reactions |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | - | $\mathrm{v}_{1(\mathrm{~K} 1)}$ | $\mathrm{v}_{2(\mathrm{k} 2)}$ | $v_{3(k 3)}$ | $\mathrm{v}_{4(\mathrm{k} 4 / 2)}$ |  |
|  | A | $?:(C \mid C)$ | ?;D |  |  |  |
| ๓. | B | ! 0 |  |  |  |  |
| - ${ }^{\text {O }}$ - | C |  | !:0 | $\tau$;(E\|F) |  |  |
| . | D |  |  |  |  |  |
| $\frac{4}{0} 0$ | E |  |  |  |  |  |
|  | F |  |  |  | $\begin{aligned} & ? ; B \\ & !; 0 \end{aligned}$ |  |

Read out the resulting CGF by rows:

$$
\begin{aligned}
& A=? v_{1(k) 1)} ;(C \mid C) \oplus ? v_{2(k 2)} ; D \\
& B=!v_{1(k))} ; 0 \\
& C=!v_{2(k 2)} ; 0 \oplus \tau_{k 3} ;(E \mid F) \\
& D=0 \\
& E=0 \\
& F=? v_{4(k 4 / 2)} ; B \oplus!v_{4(k 4 / 2)} ; 0
\end{aligned}
$$



## That CGF in SPiM

$$
\begin{aligned}
& A=? v_{1(k 1)} ;(C \mid C) \oplus ? v_{2(k 2))} ; D \\
& B=!v_{1(k 1)} ; 0 \\
& C=!v_{2(k 2)} ; 0 \oplus \tau_{k 3} ;(\mathrm{E} \mid \mathrm{F}) \\
& D=0 \\
& E=0 \\
& F=? v_{4(k 4 / 2)} ; B \oplus!v_{4(k 4 / 2) ;}: 0
\end{aligned}
$$




## directive sample 10.0 <br> directive plot $A() ; B() ; C() ; D() ; E() ; F()$

val $k 1=0.001$ new v1@k1:chan
val $k 2=0.001$ new v2@k1:chan
val $k 3=1.0$
val $k 4=0.001$ new v4@k4/2.0:chan
let $A()=\operatorname{do}$ ? $\mathrm{v} 1 ;(C() \mid C())$ or ? $\mathrm{v} 2 ; D()$
and $B()=$ ! $v 1$
and $C()=$ do !v2 or delay@k3; $(E() \mid F())$
and $D()=()$
and $E()=()$
and $F()=$ do $? v 4 ; B()$ or ! $v 4$
run 300 of $(A()|B()| C()|D()| E() \mid F())$

## ... in other words

From chemical reactions $C$ to a $C G F \operatorname{Pi}(C)$ :

$$
\begin{aligned}
\operatorname{Pi}(C)=\{(X= & \oplus\left(\left(v_{i}: X \rightarrow^{k} P\right) \in C\right) \text { of }\left(\tau_{k} ; P\right) \\
& \oplus\left(\left(v_{i}: X+Y \rightarrow^{k} P\right) \in C \text { and } Y \neq X\right) \text { of }\left(? v_{i(k i)} ; P\right) \\
& \oplus\left(\left(v_{i}: Y+X \rightarrow^{k} P\right) \in C \text { and } Y \neq X\right) \text { of }\left(!v_{i(k i)} ; 0\right) \\
& \left.\oplus\left(\left(v_{i}: X+X \rightarrow^{k} P\right) \in C\right) \text { of }\left(? v_{i(k i / 2)}: P \oplus!v_{i(k i / 2)} ; 0\right)\right) \\
& \text { s.t. } X \text { is a species in } C\}
\end{aligned}
$$

## From Processes Directly to ODEs (hard)

## "Micromodels": Continuous Time Markov Chains

- The underlying semantics of stochastic $\pi$-calculus (and stochastic interacting automata). Well established in many ways.
- Automata with rates on transitions.
- "The" correct semantics for chemistry, executable.
- Gillespie stochastic simulation algorithm
- Good description of "individual" behavior
- But does not give a good sense of "collective" properties.
- Yes one can do simulation.
- Yes one can do program analysis.
- Yes one can do modelchecking.
- But somewhat lacking in "predictive power".


## "Macromodels": Ordinary Differential Equations

- Micromodels have lots of advantages
- Compositional, compact, mechanistic, etc.
- But they always ask:
- "Yes, but how does you automata model relate to the 75 ODE models in the literature?"
- From processes/automata to ODEs directly:
- In principle: just write down the Rate Equation:

The speed of interaction is proportional to the number of possible interactions.

- Determine the set of all possible states $S$ of each process.
- Determine the rates of the transitions between such states.
- Let [S] be the "number of processes in state $S$ " as a function of time.
- Define for each state S:
$[S]^{\bullet}=($ rate of change of the number of processes in state $S$ )
Cumulative rate of transitions from any state $S^{\prime}$ to state $S$, times [ $S^{\prime}$ ], minus cumulative rate of transitions from $S$ to any state $S^{\prime \prime}$, times [S].

The rate of change of a quantity is the cumulative result of the (always changing) speed of the interactions that produce it and consume it. It is a global and dynamic property.

- Intuitive (rate = inflow minus outflow), but often clumsy to write down precisely.
- But why go directly from processes to ODE?

- If we first convert processes to chemical reactions, then we can convert to ODEs by standard means!


## From Processes to Chemistry

## Automata to Chemistry (by hand)

Automata to Chemistry

$$
\begin{array}{ll}
A+A & \begin{array}{l}
A+B \rightarrow B+B \\
B+A \rightarrow A+A
\end{array} \\
? a+B_{d} \rightarrow B+B_{d} \\
B+A_{d} \rightarrow A+A_{d}
\end{array}
$$

Doping
!a
! $b$

(B)

Doping

(Lack of) Compositionality

! $b$


6 reactions
$A+B \rightarrow A_{1}+B$
$A_{1}+B \rightarrow A_{2}+B$
$A_{2}+B \rightarrow B+B$
$B+A \rightarrow B_{1}+A$
$B_{1}+A \rightarrow B_{2}+A$
$B_{2}+A \rightarrow A+A$

6 reactions
$A+C \rightarrow A_{1}+C$
$A_{1}+C \rightarrow C+C$
$C+B \rightarrow C_{1}+B$
$C_{1}+B \rightarrow B+B$
$B+A \rightarrow B_{1}+A$
$B_{1}+A \rightarrow A+A$
+2 states
Doping
!a
! b
(Ad
(B)
+3 states
Doping
!a
!b
!c
(B)
+6 reactions??

$$
\begin{aligned}
& A+B_{d} \rightarrow A_{1}+B_{d} \\
& A_{1}+B_{d} \rightarrow A_{2}+B_{d} \\
& A_{2}+B_{d} \rightarrow B+B_{d} \\
& B+A_{d} \rightarrow B_{1}+A_{d} \\
& B_{1}+A_{d} \rightarrow B_{2}+A_{d} \\
& B_{2}+A_{d} \rightarrow A+A_{d}
\end{aligned}
$$

+6 reactions??
$A+C_{d} \rightarrow A_{1}+C_{d}$
$A_{1}+C_{d} \rightarrow C+C_{d}$
$C+B_{d} \rightarrow C_{1}+B_{d}$
$C_{1}+B_{d} \rightarrow B+B_{d}$
$B+A_{d} \rightarrow B_{1}+A_{d}$
$B_{1}+A_{d} \rightarrow A+A_{d}$

## Single Automata Models are $\mathrm{n}^{2}$ More Compact



Chemistry

$$
\begin{aligned}
& \mathrm{S}_{0}+\mathrm{S}_{0} \rightarrow \mathrm{~S}_{0}+\mathrm{S}_{1} \\
& \mathrm{~S}_{0}+\mathrm{S}_{1} \rightarrow \mathrm{~S}_{0}+\mathrm{S}_{2} \\
& \mathrm{~S}_{0}+\mathrm{S}_{2} \rightarrow \mathrm{~S}_{0}+\mathrm{S}_{0} \\
& \mathrm{~S}_{1}+\mathrm{S}_{0} \rightarrow \mathrm{~S}_{1}+\mathrm{S}_{1} \\
& \mathrm{~S}_{1}+\mathrm{S}_{1} \rightarrow \mathrm{~S}_{1}+\mathrm{S}_{2} \\
& \mathrm{~S}_{1}+\mathrm{S}_{2} \rightarrow \mathrm{~S}_{1}+\mathrm{S}_{0} \\
& \mathrm{~S}_{2}+\mathrm{S}_{0} \rightarrow \mathrm{~S}_{2}+\mathrm{S}_{1} \\
& \mathrm{~S}_{2}+\mathrm{S}_{1} \rightarrow \mathrm{~S}_{2}+\mathrm{S}_{2} \\
& \mathrm{~S}_{2}+\mathrm{S}_{2} \rightarrow \mathrm{~S}_{2}+\mathrm{S}_{0}
\end{aligned}
$$

$$
S_{0}+S_{0} \rightarrow S_{0}+S_{1} \quad S_{2}+S_{0} \rightarrow S_{2}+S_{1}
$$

$$
S_{0}+S_{1} \rightarrow S_{0}+S_{2} \quad S_{2}+S_{1} \rightarrow S_{2}+S_{2}
$$

$$
\mathrm{S}_{0}+\mathrm{S}_{2} \rightarrow \mathrm{~S}_{0}+\mathrm{S}_{3} \quad \mathrm{~S}_{2}+\mathrm{S}_{2} \rightarrow \mathrm{~S}_{2}+\mathrm{S}_{3}
$$

$$
S_{0}+S_{3} \rightarrow S_{0}+S_{0} \quad S_{2}+S_{3} \rightarrow S_{2}+S_{0}
$$

$$
S_{1}+S_{0} \rightarrow S_{1}+S_{1} \quad S_{3}+S_{0} \rightarrow S_{3}+S_{1}
$$

$$
S_{1}+S_{1} \rightarrow S_{1}+S_{2} \quad S_{3}+S_{1} \rightarrow S_{3}+S_{2}
$$

$$
S_{1}+S_{2} \rightarrow S_{1}+S_{3} \quad S_{3}+S_{2} \rightarrow S_{3}+S_{3}
$$

$$
S_{1}+S_{3} \rightarrow S_{1}+S_{0} \quad S_{3}+S_{3} \rightarrow S_{3}+S_{0}
$$

3 states
(2*3 transitions) $=$
$3^{2}$ reactions

4 states
(2*4 transitions)

42 reactions

## Process Normal Forms to Chemistry

## Chemical Parametric Form (CPF)

$$
\begin{array}{ll}
E & ::=X_{1}\left(p_{1}\right)=M_{1}, \ldots, X_{n}\left(p_{n}\right)=M_{n} \\
M & ::=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n} \\
P & ::=X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right) \\
\pi & ::=\tau_{r} \quad ? n(p) \quad!n(p) \\
C P F & ::=E, P
\end{array}
$$

Definitions $\quad(n \geq 0)$
Molecules $\quad(n \geq 0)$
Solutions
$(n \geq 0)$
Interactions
with initial conditions
$\oplus$ is stochastic choice (vs. + for chemical reactions)
0 is the null solution ( $\mathrm{P}|0=0| \mathrm{P}=\mathrm{P}$ ) and null molecule $(M \oplus 0=0 \oplus M=M)\left(\tau_{0} ; P=0\right)$
$X_{i}$ are distinct in $E, p$ are vectors of names
$p$ are vectors of distinct names when in binding position Each free name $n$ in $E$ is assigned a fixed rate $r$ :
written either $n_{(r)}$, or $\rho_{\text {CPF }}(n)=r$.

Random example

$$
\begin{aligned}
& P(a, b)=? a(c) ;(P(c, b) \mid Q(c)) \\
& Q(a)=!a(a) ; 0 \\
& P\left(x_{(r)}, Y_{(s)}\right) \mid Q\left(x_{(r)}\right)
\end{aligned}
$$

Repressilator

```
Neg(a,b) = ?a; Inh(a,b) \oplus \tau < ; (Tr(b)| Neg(a,b))
    Inh(a,b)= 㖕; Neg(a,b)
```



```
    Neg(x,y)| Neg(y,z)|Neg(z,x)
```


## CPF to Chemistry: the CGF Case

CGF = The CPF subset with no communication, no parametric definitions.

$$
\begin{array}{ll}
E::=X_{1}=M_{1}, \ldots, X_{n}=M_{n} & \text { Definitions } \\
M:(n \geq 0) \\
M:=\pi_{1} \cdot P P_{1} \oplus \ldots \oplus \pi_{n}: P_{n} & \text { Molecules } \quad(n \geq 0) \\
P::=X_{1}|\ldots| X_{n} & \text { Solutions } \quad(n \geq 0) \\
\pi::=\tau_{n} ? n_{(r)}!n_{(r)} & \text { Interactions (delay, input, output) } \\
C G F::=E, P & \text { Definitions with Initial Conditions }
\end{array}
$$

Each $X$ in $E$ is seen as a separate species.
Chemical reactions for $E: \quad$ (N.B.: $\{\ldots\}^{m}$ is a multiset, and $P$ is $P$ with all the $\mid$ changed to + )

$$
\begin{aligned}
C h_{G}(E) & :=\left\{(X \rightarrow r P) \text { s.t. }\left(X \equiv \tau_{r} ; P \oplus \ldots\right) \in E\right\}^{m} \\
& \cup^{m}\left\{(X+Y \rightarrow r P+Q) \text { s.t. } X \neq Y,\left\langle\left(X \equiv ? n_{(r)} ; P \oplus \ldots\right)\left(Y \equiv!n_{(r)} ; Q \oplus \ldots\right)\right\rangle \in E^{2}\right\} m \\
& \left.\cup^{m}\left\{\left(X+X \rightarrow r^{2 r} P+Q\right) \text { s.t. }\left(X \equiv ? n_{(r)} ; P \oplus \ldots \equiv!n_{(r)} ; Q \oplus \ldots\right)\right\rangle \in E\right\}^{m}
\end{aligned}
$$

Initial conditions for $P$ :
$\mathrm{Ch}_{G}(P):=P$

## (Note on computing the multisets)

A multiset $M \in$ Multiset(S), where $S$ is a set with equality, is a total function $S \rightarrow N a t$, which may also be written as a finite enumeration with repetitions: $\{\text {... }\}^{m}$.

Multiset binary union is the function $\cup^{m}\left(M, M^{\prime}\right)=\backslash s . M(s)+M^{\prime}(s)$. Multiset big union over a finite index set $S$ is written: $\cup^{m} i \in S$ of $M_{i}$.

- The shorthand: $\left\{\left(X \rightarrow^{r} P\right) \text { s.t. }\left(X \equiv \tau_{r} ; P \oplus \ldots\right) \in E\right\}^{m}$ means: "for each $\left(X=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n}\right) \in E$ and for each $i$ such that $\pi_{i}=\tau_{r}$, return a copy of $\left(X \rightarrow P_{i}\right)$ ".
i.e., it is defined as the following finite union of singleton multisets:

$$
\cup m\left\{\left(X=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n^{\prime}}: P_{n}\right) \in E\right\} \quad \text { of }\left(\cup m\left\{i \text { s.t. } \pi_{i}=\tau_{r}\right\} \quad \text { of }\left\{\left(X \rightarrow^{r} P_{i}\right)\right\}^{m}\right)
$$

- The shorthand: $\left\{(X+Y \rightarrow r P+Q)\right.$ s.t. $\left.X \neq Y,\left\langle\left(X \equiv ? n_{(r)} ; P \oplus \ldots\right),\left(Y \equiv!n_{(r)} ; Q \oplus \ldots\right)\right\rangle \in E^{2}\right\} m$
means: "for each ordered pair $\left(X=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n}\right),\left(Y=\sigma_{1} ; Q_{1} \oplus \ldots \oplus \sigma_{n} ; Q_{n}\right)$ from $E^{2}$ with $X \neq Y$, for each $n_{(r)}$, for each $i$ such that $\pi_{i}=? n_{(r)}$, and for each $j$ such that $\sigma_{j}=!n_{(r)}$, return a copy of the reaction $\left(X+Y \rightarrow r P_{i}+Q_{j}\right)$ "
i.e., it is defined as the following finite union of singleton multisets:

$$
\begin{aligned}
& \cup^{m}\left\langle\left(X=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n}\right),\left(Y=\rho_{1} ; Q_{1} \oplus \ldots \oplus \rho_{m} ; Q_{m}\right)\right\rangle \in E^{2} \text { with } X \neq Y \\
& \text { of }\left(\cup m\left\{\langle i, j\rangle \text { s.t. } \pi_{i}=? n_{(r)}, \rho_{j}=!n_{(r)}\right\} \quad \text { of }\left\{\left(X+Y \rightarrow P_{i}+Q_{j}\right)\right\}^{m}\right)
\end{aligned}
$$

- The shorthand: $\left\{(X+X \rightarrow 2 r P+Q)\right.$ s.t. $\left.\left.\left(X \equiv ? n_{(r)} ; P \oplus \ldots \equiv!n_{(r)} ; Q \oplus \ldots\right)\right\rangle \in E\right\}^{m}$
means: "for each $\left(X=\pi_{1} ; P_{1} \oplus \ldots, \ldots \pi_{n} ; P_{n}\right)$ in $E$, for each $n_{(r)}$, for each $i$ such that $\pi_{i}=? n_{(r)}$, and for each $j$ such that $\pi_{\mathrm{j}}=!n_{(r)}$, return a copy of the reaction $\left(X+X \rightarrow{ }^{2 r} P_{i}+P_{j}\right)$ "


## Example

$\mathrm{Ch}_{6}(\mathrm{E}, \mathrm{P})$ :
$\begin{array}{ll}\mathrm{Na}+\mathrm{Cl} \rightarrow \rightarrow^{\mathrm{p}(\mathrm{i})} & \mathrm{Na}^{+}+\mathrm{Cl}^{-} \\ \mathrm{Na}^{+}+\mathrm{Cl}^{-} \rightarrow \rightarrow^{\mathrm{p}(\mathrm{d})} & \mathrm{Na}+\mathrm{Cl}^{2}\end{array}$
$\mathrm{Na}^{+}+\mathrm{Cl}^{-} \rightarrow \mathrm{p}(\mathrm{d}) \quad \mathrm{Na}+\mathrm{Cl}$
$\mathrm{Na}+\mathrm{Na}+\mathrm{Cl}+\mathrm{Cl}$
$\rho_{\mathrm{E}}\left(\mathrm{n}_{(\mathrm{r})}\right)=r$

$$
\begin{aligned}
C h_{G}(E): & : \ldots \\
& \cup^{m}\left\{(X+Y \rightarrow r P+Q) \text { s.t. } X \neq Y,\left\langle\left(X \equiv ? n_{(r)} ; P \oplus \ldots\right),\left(Y \equiv!n_{(r)} ; Q \oplus \ldots\right)\right\rangle \in E^{2}\right\} m
\end{aligned}
$$

## Subtler Examples

These are not finite state systems, but finite species systems are ok!


Unbounded state, but only 1 species. No problem!

Multisets:
The same interaction can occur multiple times and must be taken into account:


Symmetric reactions:

$$
\begin{array}{ll}
\mathrm{E}: & C(E): \\
X=!a ; 0 \oplus ? a ; Y & X+X \rightarrow \rightarrow^{2 p(a)} Y
\end{array}
$$

The rate of a was pre-halved and must be restored.

## CPF to Chemistry: Handling Parameters

Consider first the CPF subset with no communication (pure ?n, !n).
Grounding (replace parameters with constants)

$$
\begin{aligned}
& E::=X_{1}\left(p_{1}\right)=M_{1}, \ldots, X_{n}\left(p_{n}\right)=M_{n} \\
& M::=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n}: \cdot P_{n} \\
& P::=X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right) \\
& \pi::=\tau_{r} ? n \text { ?n }
\end{aligned}
$$

where $X / p$ is a name in bijection with $\langle X, p\rangle$ (each $X / p$ is seen as a separate species)
$/\left(\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n}\right)=_{\text {def }} \pi_{1} ; /\left(P_{1}\right) \oplus \ldots \oplus \pi_{n} ; /\left(P_{n}\right)$
$/\left(X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right)\right)={ }_{\text {def }} X_{1} / p_{1}|\ldots| X_{n} / p_{n}$

Let $N$ be the set of free names occurring in $E$ (names not bound by definitions).
$E_{G}$ is the Parametric Explosion of $E$ (still a finite species system) computed by replacing parameters with all combinations of free names in $E$
$E_{G}:=\left\{(X / q=/(M\{p \leftarrow q\}))\right.$ s.t. $(X(p)=M) \in E$ and $\left.q \in N^{\# p}\right\}$
$P_{G}:=/ P \quad$ (simply ground the given initial conditions once)
$E_{G}$ is a CGF! To obtain the chemical reactions $C_{p}(E)$, just compute $C_{G}\left(E_{G}\right)$
$C h_{P}(E)=C h_{G}\left(E_{G}\right)$

## Example

E, $P$ (with free names $x, y$ ):


$$
\begin{aligned}
& Q / x=!x ; 0 \\
& Q / y=!y ; 0 \\
& P / x, y=? x ;(P / x, y \mid Q / x) \\
& P / y, x=? y ;(P / y, x \mid Q / y) \\
& P / x, x=? x ;(P / x, x \mid Q / x) \\
& P / y, y=? y ;(P / y, y \mid Q / y) \\
& P / x, y \mid Q / x
\end{aligned}
$$

$C$ :

$$
\begin{aligned}
& Q / x+P / x, y \rightarrow p(x) P / x, y+Q / x \\
& Q / x+P / x, x \rightarrow p(x) P / x, x+Q / x \\
& Q / y+P / y, x \rightarrow p(y) P / y, x+Q / y \\
& Q / y+P / y, y \rightarrow p(y) P / y, y+Q / y \\
& P / x, y+Q / x
\end{aligned}
$$

Chemical Reactions
Reduced $C$
(from initial conditions)

$$
\begin{aligned}
& Q / x+P / x, y \rightarrow p(x) P / x, y+Q / x \\
& P / x, y+Q / x
\end{aligned}
$$

Note that $E_{G}$ is not necessarily the "full" explosion (there may not be enough free variables in $E, P$ to generate it). In particular, if $E, P$ is closed (no free variables), then $E_{G}$ contains only the parameterless definitions, which must then be closed and hence unrelated to any parametric definitions. One can add initial conditions $X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right)$ where $X_{i}$ are all the definitions in $E$ and $p_{1} \ldots p_{n}$ are made of distinct (free) variables. Then $E_{G}$ becomes the most general set of chemical reactions from $E$, out of any possible initial conditions.

## Iterative Algorithm

The chemical reactions for $E$, computed directly from the parametric explosion $E_{G}$ are highly redundant because they include all the parameter permutation symmetries. The following iterative algorithm, for the parametric case, computes a subset of $E_{G}$ from the initial conditions of $E, P$. It produces a (usually) much smaller although not necessarily minimal set $C$.

## Here:

$C$ is the incrementally built set of chemical reactions
$E_{c}$ (a CGF!) is the incrementally built set of grounded definitions from the original $E$.
initialization

$$
E_{c}:=\{(X / q=/(M\{p \leftarrow q\})) \text { s.t. } X(q) \text { occurs in } P \text { and }(X(p)=M) \in E\} \quad \text { (from initial conditions } P \text { ) }
$$

iteration

$$
\begin{aligned}
& C:=C h_{G}\left(E_{C}\right) \\
& E_{c}:=E_{c} \cup\{(X / q=/(M\{p \leftarrow q\})) \text { s.t. } X / q \text { occurs in } C \text { and }(X(p)=M) \in E\}
\end{aligned}
$$

termination
if $E_{C}^{\prime}=E_{C}$ then stop and return $\left(C_{1} /(P)\right)$ else $E_{C}:=E_{C}^{\prime}$ and iterate.
The algorithm terminates because $E_{c}$ never shrinks and is always a subset of $E_{G}$, which is finite.

## Example: $\operatorname{Neg}(x, x)$

```
\(E=\)
    \(\operatorname{Neg}(a, b)=? a ; \operatorname{Inh}(a, b) \oplus \tau_{\varepsilon^{\prime}}(\operatorname{Tr}(b) \mid \operatorname{Neg}(a, b))\)
    \(\operatorname{Inh}(a, b)=\tau_{\eta} ; \operatorname{Neg}(a, b)\)
    \(\operatorname{Tr}(b)=!b ; \operatorname{Tr}(b) \oplus \tau_{\delta^{\prime}} 0\)
    \(\operatorname{Neg}(x, x)\)
----- initialization -----
\(E_{c}:=\left\{\right.\) Neg \(\left./ x, x=? x ; \operatorname{Inh} / x, x \oplus \tau_{\varepsilon^{\prime}}(\operatorname{Tr} / x \mid \operatorname{Neg} / x, x)\right\}\)
----- iteration 1 -----
\(C:=\left\{\right.\) Neg/x, \(x \rightarrow^{\varepsilon} \operatorname{Tr} / x+\) Neg/x, \(\left.x\right\}\)
\(E_{c}:=\left\{\right.\) Neg \(/ x, x=? x ; \operatorname{Inh} / x, x \oplus \tau_{\varepsilon^{\prime}}(\operatorname{Tr} / x \mid \operatorname{Neg} / x, x)\)
        \(\left.\operatorname{Tr} / \mathrm{x}=!\mathrm{x} ; \mathrm{Tr} / \mathrm{x} \oplus \tau_{\delta^{\prime}} 0\right\}\)
----- iteration 2 -----
\(C:=\left\{\right.\) Neg \(/ x, x \rightarrow \rightarrow^{\varepsilon} \operatorname{Tr} / x+\) Neg \(/ x, x\)
    \(\operatorname{Tr} / x \rightarrow{ }^{\delta} 0\)
    \(\operatorname{Tr} / x+\operatorname{Neg} / x, x \rightarrow p(x) \operatorname{Tr} / x+\operatorname{Inh} / x, x\}\)
\(E_{c}:=\left\{\mathrm{Neg} / \mathrm{x}, \mathrm{x}=? \mathrm{P} ; \operatorname{Inh} / \mathrm{x}, \mathrm{x} \oplus \tau_{\varepsilon^{\prime}}(\operatorname{Tr} / \mathrm{x} \mid \mathrm{Neg} / \mathrm{x}, \mathrm{x})\right.\)
    \(\operatorname{Tr} / x=!x ; \operatorname{Tr} / x \oplus \tau_{s^{\prime}} ; 0\)
    \(\left.\operatorname{Inh} / x, x=\tau_{n} ; \operatorname{Neg} / x, x\right\}\)
```

----- iteration 3 -----
$C:=\left\{\right.$ Neg $/ x, x \rightarrow^{\varepsilon} \operatorname{Tr} / x+$ Neg $/ x, x$
$\mathrm{Tr} / \mathrm{x} \rightarrow \mathrm{T}^{\delta} 0$
$\operatorname{Tr} / x+\operatorname{Neg} / x, x \rightarrow p(x) \operatorname{Tr} / x+\operatorname{Inh} / x, x$
Inh/ $x, x \rightarrow \eta$ Neg/ $x, x\}$
$E_{c}:=$ no change
----- termination -----

```
Neg/x,x->\mp@code{Tr}/x+Neg/x,x
Tr/x -> }
Tr/x + Neg/x,x ->p(x) Tr/x + Inh/x,x
```



```
Neg/x,x
```


## Example: $\operatorname{Tr}(x) \mid \operatorname{Neg}(x, y)$

```
\(E=\)
    \(\operatorname{Neg}(a, b)=? a ; \operatorname{Inh}(a, b) \oplus \tau_{\varepsilon^{\prime}}(\operatorname{Tr}(b) \mid \operatorname{Neg}(a, b))\)
    \(\operatorname{Inh}(a, b)=\tau_{\eta} ; \operatorname{Neg}(a, b)\)
    \(\operatorname{Tr}(\mathrm{b})=!\mathrm{b} ; \operatorname{Tr}(\mathrm{b}) \oplus \tau_{\delta^{\prime}} 0\)
    \(\operatorname{Tr}(x) \mid \operatorname{Neg}(x, y)\)
----- initialization -----
\(E^{c}:=\left\{\operatorname{Tr} / x=!x ; \operatorname{Tr} / x \oplus \tau_{\delta} ; 0\right.\)
    Neg/x,y = ?x; Inh/x,y \(\left.\oplus \tau_{\varepsilon^{\prime}}(T r / y \mid N e g / x, y)\right\}\)
----- iteration 1 -----
\(C:=\left\{\mathrm{Tr} / \mathrm{x} \rightarrow^{\delta} 0\right.\)
    Neg/x,y \(\rightarrow^{\varepsilon} \operatorname{Tr} / y+\mathrm{Neg} / x, y\)
    \(\operatorname{Tr} / x+\) Neg \(/ x, y \rightarrow p(x) \operatorname{Tr} / x+\operatorname{Inh} / x, y\}\)
\(E^{c}:=\left\{\operatorname{Tr} / x=!x ; \operatorname{Tr} / x \oplus \tau_{\delta} ; 0\right.\)
    Neg/x,y = ?x; Inh/x,y \(\oplus \tau_{\varepsilon^{\prime}}(\) Tr/y \(\mid\) Neg/x,y)
    \(\operatorname{Tr} / \mathrm{y}=\mathrm{l} \mathrm{y} ; \mathrm{Tr} / \mathrm{y} \oplus \tau_{\delta^{\prime}} 0\)
    \(\operatorname{Inh} / x, y=\tau_{\eta} ;\) Neg/x,y \}
```

----- iteration 2 -----
$C:=\left\{\operatorname{Tr} / x \rightarrow^{\delta} 0\right.$
Neg/x,y $\rightarrow^{\varepsilon}$ Tr/y + Neg/x,y
$\mathrm{Tr} / x+\mathrm{Neg} / x, y \rightarrow{ }^{p(x)} \operatorname{Tr} / x+\operatorname{Inh} / x, y$
$\operatorname{Tr} / \mathrm{y} \rightarrow{ }^{\delta} 0$
Inh/x,y $\rightarrow{ }^{n}$ Neg/x,y\}
$E^{c}:=$ no change
----- termination -----

```
\(\mathrm{Tr} / \mathrm{x} \rightarrow{ }^{\delta} 0\)
Neg/x,y \(\rightarrow^{\varepsilon}\) Tr/y + Neg/x,y
Tr/x + Neg/x,y \(\rightarrow{ }^{p(x)} \operatorname{Tr} / x+\operatorname{Inh} / x, y\)
\(\operatorname{Tr} / y \rightarrow{ }^{\delta} 0\)
    \(\operatorname{Inh} / x, y \rightarrow \eta\) Neg/x,y
    \(\mathrm{Tr} / \mathrm{x}+\mathrm{Neg} / \mathrm{x}, \mathrm{y}\)
```


## CPF to Chemistry: Handling Communication

Grounding (replace parameters with constants)
just one main change: now also convert each input parameter into a ground choice of all possible inputs
$N$ is the set of free names in $E, P$
\#p is the length of $p$

$$
\begin{aligned}
& E::=X_{1}\left(p_{1}\right)=M_{1}, \ldots, X_{n}\left(p_{n}\right)=M_{n} \\
& M::=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n} \\
& P::=X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right) \\
& \pi::=\tau_{r} ? n(p) \ln (p)
\end{aligned}
$$

$n / p$ is a name in bijection with $\langle n, p\rangle$
$X / p$ is a name in bijection with $\langle X, p\rangle$
(each $X / p$ is seen as a separate species)
$/_{N}\left(\tau_{r} ; P\right)=\tau_{r} ; /_{N}(P)$
$/ N\left(n_{(r)}(p) ; P\right)=!n / P(r) ; /_{N}(P)$
$/ N_{N}\left(? n_{(r)}(p) ; P\right)=\oplus(q \in N \# p)$ of $? n / q_{(r)} ; / N_{N}(P\{p \leftarrow q\})$
$/_{N}\left(\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n}\right)=/_{N}\left(\pi_{1} ; P_{1}\right) \oplus \ldots \oplus /_{N}\left(\pi_{n} ; P_{n}\right)$
$/_{N}\left(X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right)\right)=X_{1} / p_{1}|\ldots| X_{n} / p_{n}$
$E_{G}$ is again the Parametric Explosion of $E$

$$
\begin{aligned}
& E_{G}:=\left\{\left(X / q=/_{N}(M\{p \leftarrow q\})\right) \text { s.t. }(X(p)=M) \in E \text { and } q \in N \# p\right\} \\
& P_{G}:=/{ }_{N}(P) \quad \text { (simply ground the given initial conditions once) } \\
& C h(E)=C h_{G}\left(E_{G}\right) \quad E_{G} \text { is a again a CGF! }
\end{aligned}
$$

## Example

$E$ (with free names $x, y$ ):

$$
\begin{aligned}
& \begin{array}{l}
P(a, b)=? a(c) ;(P(c, b) \mid Q(c)) \\
Q(a)=!a(a) ; 0 \\
P(x, y) \mid Q(x)
\end{array} \\
& E_{G}:
\end{aligned}
$$

Reduced C
(from initial conditions)

$$
\begin{aligned}
& Q / x+P / x, y \rightarrow p(x) P / x, y+Q / x \\
& P / x, y+Q / x
\end{aligned}
$$

## Example

$E$ (with free names $x, y$ ):

$$
\begin{aligned}
& P(a, b)=? a(c) ;(P(c, b) \mid Q(c, c)) \\
& Q(a, b)=!a(a) ; 0 \oplus!b(b) ; 0 \\
& P(x, y) \mid Q(x, y)
\end{aligned}
$$

$E_{G}:$

$$
\begin{aligned}
& Q / x, y=!x / x ; 0 \oplus!y / y ; 0 \\
& Q / y, x=!y / y ; 0 \oplus!x / x ; 0 \\
& Q / x, x=!x / x ; 0 \oplus!x / x ; 0 \\
& Q / y, y=!y / y ; 0 \oplus!y / y ; 0 \\
& P / x, y= ? x / x ;(P / x, y \mid Q / x, x) \oplus \\
& ? x / y ;(P / y, y \mid Q / y, y) \\
& P / y, x= ? y / x ;(P / x, x \mid Q / x, x) \oplus \\
& ? y / y ;(P / y, x \mid Q / y, y) \\
& P / x, x= ? x / x ;(P / x, x \mid Q / x, x) \oplus \\
& ? x / y ;(P / y, x \mid Q / y, y) \\
& P / y, y= ? y / x ;(P / x, y \mid Q / x, x) \oplus \\
& ? y / y ;(P / y, y \mid Q / y, y) \\
& P / x, y \mid Q / x
\end{aligned}
$$

$C$ :

$$
\begin{aligned}
& Q / x, y+P / x, y \rightarrow p(x) P / x, y+Q / x, x \\
& Q / x, y+P / x, x \rightarrow p(x) P / x, x+Q / x, x \\
& Q / x, y+P / y, x \rightarrow p(y) \quad P / y, x+Q / y, y \\
& Q / x, y+P / y, y \rightarrow p(y) \quad P / y, y+Q / y, y \\
& Q / y, x+P / y, x \rightarrow p(y) \quad P / y, x+Q / y, y \\
& Q / y, x+P / y, y \rightarrow p(y) \quad P / y, y+Q / y, y \\
& Q / y, x+P / x, y \rightarrow p(x) P / x, y+Q / x, x \\
& Q / y, x+P / x, x \rightarrow p(x) P / x, x+Q / x, x \\
& Q / x, x+P / x, y \rightarrow{ }^{2 p(x) P / x, y+Q / x, x} \\
& Q / x, x+P / x, x \rightarrow 2 p(x) P / x, x+Q / x, x \\
& Q / y, y+P / y, x \rightarrow \rightarrow^{2 \rho(y) P / y, x+Q / y, y} \\
& Q / y, y+P / y, y \rightarrow{ }^{2 \rho(y) P / y, y+Q / y, y} \\
& P / x, y+Q / x, y
\end{aligned}
$$

reduced $C$ (from initial conditions)

$$
\begin{aligned}
& Q / x, y+P / x, y \rightarrow \rightarrow^{p(x)} P / x, y+Q / x, x \\
& Q / x, x+P / x, y \rightarrow \rightarrow^{2 p(x) P / x, y+Q / x, x} \\
& P / x, y+Q / x, y
\end{aligned}
$$

## Exponential Explosion via Communication

Parametric form (size 3).

green $=$ free names with associated rates

Exponentially larger ground form (size 8).

$$
\begin{aligned}
& X / t_{1}, f_{1}=!t_{1} / f_{1} ; X / f_{1}, t_{1} \\
& X / f_{1}, t_{1}=!f_{1} / t_{1} ; X / t_{1}, f_{1} \\
& X / t_{2}, f_{2}=!t_{2} / f_{2} ; X / f_{2}, t_{2} \\
& X / f_{2}, t_{2}=!f_{2} / t_{2} ; X / t_{2}, f_{2} \\
& X / t_{3}, f_{3}=!t_{3} / f_{3} ; X / f_{3}, t_{3} \\
& X / f_{3}, t_{3}=!f_{3} / t_{3} ; X / t_{3}, f_{3}
\end{aligned}
$$

$$
A / t_{1}, t_{2}, t_{3}=
$$

$$
? t_{1} / f_{1} ; A / f_{1}, t_{2}, t_{3} \oplus
$$

$$
? t_{2} / f_{2}, A / t_{1}, f_{2}, t_{3} \oplus
$$

$$
? t_{3} / f_{3} ; A / t_{1}, t_{2}, f_{3}
$$

$$
\mathrm{A} / \mathrm{f}_{1}, \mathrm{t}_{2}, \mathrm{t}_{3}=
$$

$$
? f_{1} / t_{1} ; A / t_{1}, t_{2}, t_{3} \oplus
$$

$$
? t_{2} / f_{2} ; A / t_{1}, f_{2}, t_{3} \oplus
$$

$$
? t_{3} / f_{3} ; A / t_{1}, t_{2}, f_{3}
$$

...etc for the other bit flips

$$
A / t_{1}, t_{2}, t_{3}
$$

$$
X / t_{1}, f_{1}\left|X / t_{2}, f_{2}\right| X / t_{3}, f_{3}
$$

We have not fully expanded the input summations in $A(\ldots)$ : only enough for the channels that are actually used: e.g., there would never be a communication on channel $t_{1} / f_{2}$.

## CPF Iterative Algorithm

The same as before, just using the new grounding $/ N$ (based on the set $N$ of free names of $E, P$, for expanding inputs).
initialization

$$
N=f n(E, P)
$$

$$
E_{C}:=\left\{\left(X / q=/_{N}(M\{p \leftarrow q\})\right) \text { s.t. } X(q) \text { occurs in } P \text { and }(X(p)=M) \in E\right\} \quad \text { (initial conditions) }
$$

iteration

$$
C:=\quad h_{G}\left(E_{C}\right)
$$

$$
E_{c}^{\prime}:=E_{c} \cup\left\{\left(X / q=/_{N}(M\{p \leftarrow q\})\right) \text { s.t. } X / q \text { occurs in } C \text { and }(X(p)=M) \in E\right\}
$$

termination
if $E_{C}^{\prime}=E_{C}$ then stop and return $\left(C, /_{N}(P)\right)$ else $E_{C}:=E_{C}$ and iterate.

## A General $\pi$-calculus Normal Form

$$
\begin{array}{ll}
\mathrm{E} & ::=X_{1}\left(\mathbf{p}_{1}\right)=\left(v \mathbf{q}_{1}\right) M_{1}, \ldots, X_{n}\left(p_{n}\right)=\left(v q_{n}\right) M_{n} \\
M & ::=\pi_{1} ; P_{1} \oplus \ldots \oplus \pi_{n} ; P_{n} \\
P & ::=X_{1}\left(p_{1}\right)|\ldots| X_{n}\left(p_{n}\right) \\
\pi & ::=\tau_{r} \quad ? n(p) \ln (p) \\
N F & ::=E, P
\end{array}
$$

Definitions $\quad(n \geq 0)$

Molecules $\quad(n \geq 0)$
Solutions $\quad(n \geq 0)$
Interactions
with initial conditions

Any $\pi$-calculus process can be written in this normal form.
But not all of those can be translated to chemical reactions, because we need to restrict to a finite number of species.

Example (complexation):

$$
\begin{aligned}
& P(a)=\left(v n_{(d)}\right)!a(n) ; P^{\prime}(a, n) \\
& P^{\prime}(a, n)=!n ; P(a) \\
& Q(a)=? a(n) ; Q^{\prime}(a, n) \\
& Q^{\prime}(a, n)=? n ; Q(a) \\
& P\left(x_{(c)}\right) \mid Q\left(x_{(c)}\right)
\end{aligned}
$$

Example (enzymatic reaction):

$$
\begin{aligned}
& E(a)=\left(v n_{(d)}, m_{(p)}\right)!a(n, m) ; E^{\prime}(a, n, m) \\
& E^{\prime}(a, n, m)=!n ; E(a) \oplus!m ; E(a) \\
& S(a)=? a(n, m) ; S^{\prime}(a, n, m) \\
& S^{\prime}(a, n, m)=? n ; S(a) \oplus ? m ; P() \\
& P()=\tau_{0} ; 0 \\
& E\left(x_{(c)}\right)\left|S\left(x_{(c)}\right)\right| S\left(x_{(c)}\right)
\end{aligned}
$$

## Summary

- From Chemistry to ODEs
- Via the stoichiometric matrix
- From Chemistry to Processes
- Via the interaction matrix
- From Processes to Chemistry
- Ground processes: by analysis of normal form interactions
- Parametric processes: by parametric explosion
- Compositionality
- Processes are compositional in components (you just "add components") and chemistry is compositional in reactions (you just "add reactions"). But:
- Adding reactions has a linear effect on the number of components (adding one reaction adds one or two more process actions).
- Adding components has a quadratic effect on the number of reactions (adding one species may add one reaction for each species).
- ODEs grow quadratically along with the chemical reactions (stoichiometric matrix has one column per reaction).
- Parameterization is available only on processes.

> Q?

