

Motivation

- Give substance to the claim that "cells compute"
 - Yes, but *what* do they compute?
- Catch nature red-handed in the act of running a computational task
 - Something that a computer scientist would recognize as an *algorithm*





A Biological Implementation

Approximate Majority (AM)



Bistable Even when x=y (stochastically)

Fast O(log n) convergence time

Robust to perturbation above a threshold, initial majority wins *whp*

Dana Angluin - James Aspnes - David Eisenstat A Simple Population Protocol for Fast Robust Approximate Majority 2007

Epigenetic Switch



Figure 1. Basic Ingredients of the Model



Motivation (cont'd)

- We can claim that the epigenetic switch is a *direct* biological implementation of an algorithm
 - Although we may have to qualify that with some notion of approximation of the (enzymatic) kinetics
- In most cases the biological implementation seems more *indirect* or *obfuscated*
 - "Nature is subtle but not malicious Einstein" Ha! think again!
 - Other implementations of Approximate Majority seem convoluted and... approximate



Obfuscated Implementations

• GW is a better cell cycle switch than [the traditional switch]



• GW is how the cell cycle switch "really works"



Motivation (cont'd)

- When does a biologically messy network X "implement" some ideal algorithm Y?
 - \cdot Pushed coauthors into thinking about approximate stochastic bisimulation metrics for CTMCs
 - But they didn't come back...
- · Some networks behave similarly because "their ODEs are just equivalent" [David S.]
 - When are CRNs "deterministically equivalent"?
 - Or better, when do trajectories of one CRN "collapse" into trajectories of another?
 - This can be answered on the *static structure* of CRNs as opposed to their kinetics.
 - Independently on rates and initial conditions (of one of the two networks).













Network Emulation: NCC emulates MI

 For any rates and initial conditions of MI we can find some rates and initial conditions of NCC such that the (18) trajectories of NCC retrace those (6) of MI









Emulations Compose: NCC emulates AM

• The (18) trajectories NCC can *always* retrace those (3) of AM



Emulation in Context



AM-AM Oscillator



AM-MI Oscillator





 $m \in MI \rightarrow AM$ is an emulation: it maps $z \rightarrow x$ and $\sim w \rightarrow x$

We can replace AM with MI in a context. The mapping m tells us how to wire MI to obtain an overall emulation:

Each influence crossing the dashed lines into x is replaced by a similar influence into both z and $\sim w$. The latter is the same as an opposite influence into w (shown).

Each influence crossing the dashed lines out of x is replaced by a similar influence from the same side of *either z or* $\sim w$. The latter is the same as a similar influence from the opposite side of w (shown), and the same as an opposite influence from the same side of w.







A Theory of Network Emulation (with thanks to David Soloveichik)

- So far, evidence is empirical
 - · Specific simulations based on a choice of parameters
- But indeed...
 - We can show analytically that, GW, NCC, etc. are *exactly* and *always* as good as AM
 - Where *exactly* means *numerically* as good, not just in the same complexity class
 - And *always* means for *any* choice of rates and initial conditions (of the target network)
- A network *emulates* another network:
 - When it can *exactly* reproduce the kinetics of another network for *any* choice of rates and initial conditions (of the other network)
 - We aim to show that e.g. the cell cycle switch can emulate AM in that sense
 - And moreover that the emulation is algorithmic: it is determined by static network *structure* (including rate constants and stoichiometric constants), not by random kinetics

When can a Network Emulate Another?

- What kind of morphisms guarantee emulation?
 - do they preserve network structure?
 - do they preserve stoichiometry?



Chemical Reaction Networks

- A CRN is a pair (S, R) where
 - $S = \{s_1, ..., s_n\}$ a finite set of species • $R = \{r_1, ..., r_m\}$ a finite set of reactions^(*)

• Reactions
$$r = \rho \rightarrow^k \pi \in R$$

with complexes $\rho, \pi \in \mathbb{N}^S$
stoichiometric numbers ρ_s, π_s for $s \in S$
and rate constants $k > 0$

• The stoichiometry of s in $\rho \rightarrow^k \pi$ is:

$$\eta(s, \rho \to^k \pi) = \pi_s - \rho_s$$
$$\varphi(s, \rho \to^k \pi) = k \cdot (\pi_s - \rho_s)$$

$$S = \{A, B, C\}$$
$$R = \{r\}$$

$$r = 2A + B \to^{k} A + 3C$$

$$\rho_{A} = 2, \ \rho_{B} = 1, \ \rho_{C} = 0$$

$$\pi_{A} = 1, \ \pi_{B} = 0, \ \pi_{C} = 3$$



$$\eta(A,r) = -1$$
 net stoichiometry
 $\varphi(A,r) = -k$ (instantaneous) stoichiometry

$${}^{(*)}\rho \to^k \pi, \rho \to^{k'} \pi \in R \quad \Rightarrow \quad k = k'$$

CRN Morphisms

A CRN morphism from (S, R) to (\hat{S}, \hat{R}) written $m \in (S, R) \rightarrow (\hat{S}, \hat{R})$

is a pair of maps $m = (m_S, m_R)$ a species map $m_S \in S \rightarrow \hat{S}$ a reaction map $m_R \in R \rightarrow \hat{R}$

extended to a complex map $m_{\mathcal{S}} \in \mathbb{N}^{S} \to \mathbb{N}^{\hat{S}}$ linearly: $m_{\mathcal{S}}(\rho)_{\hat{S}} = \sum_{s \in m_{\mathcal{S}}^{-1}(\hat{s})} \rho_{s}$

(sometimes omitting the subscripts on m)

Mappings (symmetries) between two networks



3 Key Morphisms

- A morphism $m \in (S, R) \rightarrow (\hat{S}, \hat{R})$ is
 - a CRN homomorphism if $m_{\mathcal{R}}$ is determined by $m_{\mathcal{S}}$:

 $m_{\mathcal{R}}(\rho \to^k \pi) = m_{\mathcal{S}}(\rho) \to^k m_{\mathcal{S}}(\pi)$

• a *CRN reactant morphism* if $m_{\mathcal{R}}$ is determined by $m_{\mathcal{S}}$ on reactants. $\exists \hat{k}, \hat{\pi}$:

• a CRN stoichiomorphism if:

 $\varphi, \widehat{\varphi}$ are the respective stoichiometric matrices $\rho, \widehat{\rho}$ are the respective reactant matrices $m_{\mathcal{S}}, m_{\mathcal{R}}$ are the characteristic 0-1 matrices of $m_{\mathcal{S}}, m_{\mathcal{R}}$ $m_{\mathcal{S}}(s, \widehat{s}) = 1$ if $m_{\mathcal{S}}(s) = \widehat{s}$ else 0

$$\boldsymbol{m}_{\mathcal{S}}^{\mathrm{T}}\cdot\boldsymbol{arphi}=\widehat{\boldsymbol{arphi}}\cdot\boldsymbol{m}_{\mathcal{R}}^{\mathrm{T}}$$

$$\omega \cdot m_{\mathcal{D}} = m_{\mathcal{S}} \cdot \widehat{\omega}$$

def.

 $m_{\rm s}^{\rm T} \cdot \rho = \hat{\rho} \cdot m_{\rm p}^{\rm T}$



CRN Kinetics

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A state of a CRN (S, R) is a v \in \mathbb{R}^{S}_{+}
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The differential system of a CRN (S, R), $F \in \mathbb{R}^S_+ \to \mathbb{R}^S$



Given by the law of mass action:

$$F(\boldsymbol{\nu})(s) = \Sigma_{r \in R} \ \varphi(s, r) \cdot [r]_{\boldsymbol{\nu}}$$

Usually written as a system of coupled concentration ODEs, integrated over time: $\frac{dv_s}{dt} = F(v)(s)$

a vector of concentrations for each species

F(v)(s) gives the instantaneous change of concentration of a species in a given state

sum over all reactions of the stoichiometry of the species in the reaction times the mass action of the reaction in the state

the mass action of a reaction in state is the product of reagent concentrations according to their stoichiometric numbers:

$$[\rho \to^k \pi]_{\boldsymbol{v}} = \boldsymbol{v}^{\rho} = \Pi_{s \in S} \, \boldsymbol{v}_s^{\rho_s} \qquad _{29}$$

Kinetic Emulation

A morphism $m \in (S, R) \rightarrow (\hat{S}, \hat{R})$ is a *CRN emulation* if for the respective differential systems $F, \hat{F}, \forall \hat{v} \in \mathbb{R}^{\hat{S}}_+$:

 $F(\widehat{\boldsymbol{v}} \circ m) = \widehat{F}(\widehat{\boldsymbol{v}}) \circ m$



That is: $\forall s \in S$. $F(\hat{v} \circ m)(s) = \hat{F}(\hat{v})(m(s))$





if the derivative of s (in state $\widehat{v} \circ m$) equals the derivative of m(s) (in state \widehat{v})

if we *start* the two systems in states $\boldsymbol{v} = \boldsymbol{\hat{v}} \circ \boldsymbol{m}$ (which is a *copy* of $\boldsymbol{\hat{v}}$ according to \boldsymbol{m}) and $\boldsymbol{\hat{v}}$ resp., for each \boldsymbol{s} the solutions are equal and the derivatives are equal, hence they will have identical trajectories by determinism



Change of Rates Theorem

A change of rates for (S, R) is morphism $\iota \in (S, R) \rightarrow (S, R')$ such that $\iota(S)$ is the identity and $\iota(\rho, \pi, k) = (\rho, \pi, k')$.

Theorem: If $m \in (S, R) \to (\hat{S}, \hat{R})$ is a stoichiomorphism, then for *any* change of rates $\hat{\iota}$ of (\hat{S}, \hat{R}) there is a change of rates ι of (S, R) such that $\hat{\iota} \circ m \circ \iota^{-1}$ is a stoichiomorphism.

In fact, ι changes rates by the ratio with which $\hat{\iota}$ changes rates: $\iota(\rho, \pi, k) = \left(\rho, \pi, k \cdot \frac{\hat{k}'}{\hat{k}}\right)$ where $m(\rho, \pi, k) = (\hat{\rho}, \hat{\pi}, \hat{k})$ and $\hat{\iota}(\hat{\rho}, \hat{\pi}, \hat{k}) = (\hat{\rho}, \hat{\pi}, \hat{k}')$.

 $\begin{array}{c|c} R' & R & & \hat{R} & \hat{R}' \\ \hline (\rho',\pi',k') & & & & & \\ (\rho',\pi',k') & & & & & \\ (m^{-1}(\hat{\rho}^{-1}(\hat{\rho})\hat{\pi},\hat{k}))) & & & & & \\ m^{-1}(\hat{\rho},\hat{\pi},\hat{k}) & & & & & \\ \end{array}$

a morphism that modifies rates only

thus, for *any rates* of (\hat{S}, \hat{R}) we can match trajectories

Corollaries

- By checking only static network and morphism properties we can learn that:
 - All these networks are (at least) bistable
 - (We do not have to reanalyze the steady states of all these dynamical systems)
 - All these networks can perform *exactly* as fast as AM
 - (We do not have to reprove the complexity bounds for all these networks)





Network Emulation Morphisms

- What guarantees emulation?
 - Reactant morphism + stoichiomorphism: static, state-independent (structural) conditions
- How do you find them?
 - Emulation Theorem => they do not depend on initial conditions
 - Change of Rates Theorem => can look for rate-1 morphisms
 - E.g. test all possible rate-1 homomorphism between two networks to see if they are stoichiomorphisms

• How common are they?

• Likely relatively rare, but still many useful ones => richness of networks space

• How useful are they?

- Establish structural, algorithmic, (non-accidental) reasons for kinetic similarity
- Explain simple behavior "facets" of complicated networks
- Investigate evolutionary paths (maybe)

• How brittle are they?

- Will a perturbed trajectory of the source network converge to a trajectory of the target network?
- What about other reaction kinetics?
- What about stochastic?
 - Is there a CME Emulation Theorem?

Population Majority			
2004: Computation in networks of passively mobile finite-state sensors. Dana Angluin, James Aspnes, Zoë Diamadi, Michael J. Fischer, René Peraita. PODC'04.	Majority. The value of the majority function is 1 if there are more 1's than 0's in the input; otherwise, it is 0. The states of our protocol consist of a live bit and a counter with values in the set $\{-1, 0, 1\}$. Initially, the live	Exact Majority - 6-state Nondeterministic . (population protocol)	
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2012: The Cell Cycle Switch Computes Approximate Majority . Luca Cardeli, Attila Csikász-Nagy. Scientific Reports.		The biological cell cycle switch is a (non-obvious) implementation of approximate majority. Simulations.	
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