Reachability Computation for Switching Diffusions: Finite Abstractions with Certifiable and Tuneable Precision

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ABSTRACT

We consider continuous time stochastic hybrid systems with no resets and continuous dynamics described by linear stochastic differential equations – models also known as switching diffusions. We show that for this class of models reachability (and dually, safety) properties can be studied on an abstraction defined in terms of a discrete time and finite space Markov chain (DTMC), with provable error bounds. The technical contribution of the paper is a characterization of the uniform convergence of the time discretization of such stochastic processes with respect to safety properties. This allows us to newly provide a complete and sound numerical procedure for reachability and safety computation over switching diffusions.

Keywords

Switching diffusions; stochastic hybrid models; reachability and safety analysis; finite abstractions; time and space discretisation; numerical computations

1. INTRODUCTION

Hybrid models are natural in the context of cyber-physical systems applications, where continuous dynamics of physical variables are interleaved with discrete updates of finite-state models. In many engineering and natural systems, noise or uncertainty structured via probabilistic laws are relevant, which leads to stochastic models. In this context stochastic hybrid models encompass all these features, and their properties have been recently investigated [17, 10, 18, 3].

In this work we consider switching diffusions [27, 38, 6], models that are characterised by dynamics over a hybrid state space: continuous-time flows are determined by the solution of a mode-dependent linear diffusion process, whereas mode updates (over finitely many locations) hinge on events triggered by Poisson processes, with rates that depend on the continuous variables. As such, switching diffusions can be regarded as special instances of stochastic hybrid models, the latter dealing also with probabilistic resets between discrete-mode communications. The models considered in this work are fully observable and not subject to any form of nondeterminism (such as control inputs, as discussed in [27, 6]).

This paper investigates the problem of reachability analysis for switching diffusions, a central problem due to the duality between reachability and safety problems, and its role in the verification of many other specifications (thanks to product constructions). Whilst this is a widely investigated problem, contributions in the literature have been limited to the characterisation of this problem, with computational aspects that have been relegated to the use of approximation techniques often resorting to state-space gridding with no guarantees.

Contribution

This work provides a formal computational procedure for the reachability analysis problem over switching diffusions. As such, we address an open problem also for the special case of linear stochastic differential equations. More precisely, we provide approximation algorithms with certificates on their precision, which reduce the problem to the computation over a finite-state Markov chain. In other words, we show that probabilistic reachability can be formally computed over finite abstractions, obtained by discretising the continuous components of the models (time and space).

Related work

Stochastic hybrid models (SHS) are extensively discussed in [18, 10], and switching diffusions investigated in [27, 38, 1, 6]. The characterisation of probabilistic reachability for SHS is elaborated in [13] by means of a number of techniques, but not under the lens of computations. [13] leverages and extends theory developed for piecewise deterministic Markov processes in [20]. Further, [29] has characterised probabilistic reachability for SHS as a solution of a PDE (HJI partial differential equation), but only provided weak convergence results for its computation, based on the approximation the-
ory in [30]. A similar approach has been pursued in [34], but again without a numerical scheme with certifiable errors. It appears that the application of numerical schemes for time discretisation of SDE [28] are not helpful. [22] has extended the characterisation to constrained reachability problems. In [11, 16, 12] numerical algorithms for verification of linear SDE, obtained for Markov population processes in the limit of high population, have been given with just weak convergence results.

For discrete-time stochastic hybrid models, probabilistic reachability (and safety) have been fully characterised [3], connected with verification procedures [2, 37], formally computed via software tools [23] leveraging finite abstractions, and indeed extended to general specifications [36].

An alternative approach towards formal, finite approximations of continuous-time stochastic models is discussed in [40] and extended in [39] to switching diffusions. Noteworthy are also techniques and tools for verification of related probabilistic models based on abstractions [41], measurability conditions [24], and SMT technology [25] approaches. These techniques, alongside that of this work, are clearly distinct from statistical model checking approaches [15].

2. STOCHASTIC HYBRID PROCESSES

We consider the following class of continuous time stochastic hybrid systems with no guards or resets, which are also commonly denoted as switching diffusions. We refer the reader to [10, 18] for technical details on the measure theoretical aspects underlying these processes.

Definition 1. A switching diffusion $\mathcal{H}$ is a tuple $\mathcal{H} = (Q, K, F, G, W, \Lambda)$, where

- $Q = \{q_1, \ldots, q_P\}$ is the set of discrete modes
- $K \subseteq \mathbb{R}^m$, for $m > 0$, is the state space of the continuous dynamics. The hybrid state space is defined as $\mathcal{D} = \cup_{q \in Q} \{q\} \times K$
- $F : Q \to \mathbb{R}^{m \times m}$ is the drift term for the continuous dynamics
- $G : Q \to \mathbb{R}^{m \times q}$ is the diffusion associated to the continuous dynamics
- $W$ is a $q$-dimensional Wiener process
- $\Lambda : \mathcal{D} \times Q \to \mathbb{R}_{\geq 0}$ is an intensity function, where for $(q_i, x) \in \mathcal{D}, q_j \in Q$, we define $\Lambda((q_i, x), q_j) = \lambda_{i,j}(x)$

Let $W$ be defined in the probability space $(\Omega, \mathcal{F}, P)$ with filtration $\mathcal{F}_t$, where a filtration is a family of $\sigma$-algebras representing the information available at time $t$. Then, given $\mathcal{H}$ and an initial condition $y_0 = (x_0, q_0) \in \mathcal{D}$, the stochastic process $X = (X, \alpha)$, defined on the hybrid state space $\mathcal{D} = \cup_{q \in Q} \{q\} \times K$ is a solution of $\mathcal{H}$ if it satisfies

$$dX(t) = F(\alpha(t)) \cdot X(t)dt + G(\alpha(t)) \cdot dW(t),$$

and for $i \neq j$

$$P \left( \alpha(t + \Delta t) = q_j \mid Y(t) = (q_i, x) \right) = \lambda_{i,j}(x) \Delta t + o(\Delta t)$$

with $(X(0), \alpha(0)) = (x_0, q_0)$.

The discrete dynamics of $Y$, described by the variable $\alpha$, evolves as a jump process over the discrete state space $Q$, with jump rate dependent on the continuous part. The continuous dynamics of $Y$ evolves according to a linear diffusion. That is, when the discrete system is in a particular state, $X$ evolves according to a linear SDE driven by a Wiener process. Then, when the discrete system hits a change in its state, $X$ continues to evolve according to a different SDE without resetting its state.

Assumption 1. We introduce the following assumptions, which are standard in the literature [33, 38]:

- $\lambda_{i,j}(x)$ is a bounded and locally Lipschitz continuous function in $x$, for all $q_i, q_j \in Q$
- $|F(q)x| + |G(q)| \leq C(1 + |q|)$ for all $q \in Q$, for some constant $C$ where $|G(q)| = \sum_{i,j} |G(q)(t,j)|$
- $|F(q)x - F(q)x'| \leq D|x - x'|$ for all $q \in Q$, for some constant $D$

The first condition guarantees that, over any finite time interval, $\alpha$ almost surely jumps only a finite number of times, thus excluding Zeno behaviours. The second and third conditions guarantee that the continuous solution $X$ exists and is unique, and that it remains bounded over a finite time interval [33].

Example 1. Consider the stochastic process $X$ described by the following SDE

$$dX(t) = F \cdot X(t)dt + G \cdot dW(t)$$

with initial condition $X(0) = x_0 \in \mathbb{R}^n$. That is, $X$ is the solution of a hybrid process $\mathcal{H}$ with a singleton discrete state space $Q = \{q\}$. It is well known that the evolution of the probability distribution of the solution of a SDE over time satisfies the following Fokker-Planck equation [26]:

$$\frac{\partial p(x, t)}{\partial t} = -\sum_{i=1}^{X^q} \frac{\partial}{\partial x_i} (F(x) \cdot p(x, t)) + \frac{1}{2} \sum_{i=1}^{X^q} \sum_{j=1}^{X^q} \frac{\partial^2}{\partial x_i \partial x_j} D_{ij} p(x, t),$$

with diffusion tensor $D_{ij} = \sum_{k=1}^{X^q} G_{ik} G_{jk}$.

The following lemma guarantees that $X$, process solution of Equation 3, is a Gaussian process.

Lemma 1. [7, 35] Let $X(0)$ be a normally distributed random variable with expected value $E[X(0)] = E_{x_0}$ and covariance matrix $C_x(0) = E[X(0)X(0)^T] = C_{x_0}$. Then, $X$, as defined in (3), is a Gaussian Markov process with expected value and covariance matrix given by

$$\begin{align*}
\frac{dE[X(t)]}{dt} &= FE[X(t)] \\
E[X(0)] &= E_{x_0},
\end{align*}$$

$$\begin{align*}
\frac{dC_X(t)}{dt} &= FC_X(t) + C_X(t)^T F + G(G^T) \\
C_X(0) &= C_{x_0}.
\end{align*}$$

Lemma 1 allows us to derive the analytical solution for the expectation and variance of the solution of a linear SDE as

$$E[X(t)] = e^{Ft} E_{x_0},$$

$$C_X(t) = e^{Ft} C_{x_0} e^{F^T t} + \int_0^t e^{F(s-t)} GG^T e^{F^T (t-s)} ds.$$
3. PROBLEM DEFINITION

Given a stochastic process $Y$ with state space $D$, a target set $S \subseteq D$, which is assumed to be measurable, and a time interval $I \subseteq \mathbb{R}_{\geq 0}$, the reachability problem is defined as the search for the characterisation and computation of the probability that $Y$ will reach $S$ during $I$ from any point in $D$. This problem is dual to the safety problem, that is, computing the probability that the system will remain in a given, measurable safe region, over a given time interval. The characterisation of the two problems is thus interchangeable [2]. Reachability analysis is one of the fundamental problems in the quantitative analysis of models, and it is likewise key for the analysis of stochastic hybrid processes [14]. Model checking of Continuous Stochastic Logic (CSL) [8] reduces to computing reachability problems. Similarly, for discrete-time stochastic hybrid systems, reachability and safety play a pivotal role for model checking PCTL-like formulae [31] and more complex properties via the product construction [36].

**Problem 1. (Probabilistic Reachability)** Let $\mathcal{H}$ be a hybrid process, and $Y = (X, \alpha)$ its solution with state space $D$. Let $S \subseteq D$ be a measurable set and $I = [t_1, t_2]$ a time interval. The reachability probability for $Y$ to reach $S$ in $I$ is defined as

$$P_{\text{reach}}(Y, S, I) = \text{Prob}\{\exists t \in I \text{ s.t. } Y(t) \in S\}.$$ (6)

The safety problem is introduced as

$$P_{\text{safety}}(Y, S, I) = \text{Prob}\{\forall t \in I, Y(t) \in S\}$$

and is the dual of the reachability problem, namely

$$P_{\text{safety}}(Y, S, I) = 1 - P_{\text{reach}}(Y, S^c, I)$$

where $S^c$ is the complement of set $S$.

Analytic solutions of Problem 1 for the class of hybrid systems we consider are in general infeasible, as they would be tantamount to viscosity solutions of systems of Hamilton-Jacobi-Bellman equations [29]. In this work we instead introduce a numerical algorithm that employs time- and space discretization to solve Problem 1 – in particular the time discretisation part of the scheme is new. We further show that for the class of processes considered in this paper, the safety value computed on the discrete time and finite space Markov chain (DTMC) abstraction, obtained from the overall procedure, converges uniformly to the safety value associated to the given (continuous) switching diffusion as the discretisation parameters become zero. We also offer explicit error bounds quantifying this approximation level.

In the following illustrative example, we consider a simple SDE model, for which analytical solutions to the probabilistic safety problem exist.

**Example 2. Consider the stochastic process $X$ described by the SDE $dX(t) = GdW(t)$, where $G \in \mathbb{R}_{\geq 0}$ and $W$ is a uni-dimensional Brownian motion. Assume there is no discrete switching (and a single discrete location), then the probability density function of process $X$ is described by the following diffusion equation**

$$\frac{\partial p(x, t \mid x_0)}{\partial t} = \frac{G^2}{2} \frac{\partial^2 p(x, t \mid x_0)}{\partial x^2},$$

with initial condition $p(x, 0 \mid x_0) = \delta(x - x_0)$. For $\bar{x} \in \mathbb{R}$, we consider the safe set $S_{\bar{x}} = \{x \in \mathbb{R} : x < \bar{x}\}$. In order to solve $P_{\text{safety}}(X, S_{\bar{x}}, [0, 1])$, we need to integrate $p(x, t \mid x_0)$ with the boundary condition $P(\bar{x}, t) = 0$: this leads to the following density function for, $x < \bar{x}$,

$$p(x, t \mid x_0, \bar{x}) = \frac{1}{\sqrt{2\pi Gt}} e^{-\frac{(x-x_0)^2}{2Gt}} - e^{-\frac{(x-\bar{x}-x_0)^2}{2Gt}}.$$ 

We then obtain

$$P_{\text{safety}}(X, S_{\bar{x}}, [0, 1]) = \int p(x, 1 \mid x_0, \bar{x})dx = \text{erf} \frac{\bar{x} - x_0}{\sqrt{2G}},$$

where erf() is the Gaussian error function.

4. TIME DISCRETIZATION

Given a hybrid system $\mathcal{H}$, its solution, $Y = (X, \alpha)$, is a continuous time Markov process defined on the hybrid space $D = \cup_{q \in \Omega} \{q\} \times K$, where $K \subseteq \mathbb{R}^m$, $m > 0$. By sampling $Y$ with a fixed interval $h > 0$, we obtain a discrete-time Markov process $\tilde{Y} = (\tilde{X}, \tilde{\alpha})$ defined on the same hybrid state space $D$ and such that $\tilde{Y}(k) = Y(hk, k), k \in \mathbb{N}$.

**Definition 2.** For $k \in \mathbb{N}$, the discrete-time Markov process (DTMP) $(\tilde{Y}(k) = (\tilde{X}(k), \tilde{\alpha}(k))$ is a time homogeneous hybrid model, uniquely defined by a quadruple $(D, \sigma, T^c, T^d)$, where $(D, \sigma)$ is the measurable space inherited from $\mathcal{H}$; $T^c : A \times D \to [0, 1]$, for $A \subseteq \mathbb{R}^m$, is a continuous transition kernel; and $T^d : Q \times D \to [0, 1]$ is a discrete transition kernel. $T^c$ and $T^d$ describe the probability that the continuous and discrete components of the process transition onto a measurable set at the next discrete step, given the current state of the process. More precisely, for state $(q, x) \in D$ and Borel-measurable set $(q', A) \subseteq D$, we have that

$$T^c(A, x, q) = \text{Prob}(\tilde{X}(k + 1) \in A\mid \tilde{X}(k) = x, \tilde{\alpha}(k) = q)$$

$$T^d(q', x, q) = \text{Prob}(\tilde{\alpha}(k + 1) = q'\mid \tilde{X}(k) = x, \tilde{\alpha}(k) = q).$$

$T^c$ and $T^d$ fully characterize $\tilde{Y} = (\tilde{X}, \tilde{\alpha})$. In the following proposition we derive an analytical form for such kernels. To simplify the presentation, for the following theorem only we make a further restriction that the jump rates do not depend on the continuous state $x \in K$: $\lambda_j(x) = \lambda_j$. This assumption allows us to have a simpler form of the kernel. In order to deal with more general rate functions, we can assume that they are piecewise constant in each considered interval of time, and fix the value of $\lambda_j$ at the initial state for each time interval. As rate functions are locally Lipschitz, the distance between the true rate and the obtained $\lambda_j$ will be bounded by a term of order $O(h)$, whose error contribution can be lifted to the kernel level.

**Theorem 1.** Let $\mathcal{H} = (Q, K, F, G, W, \Lambda)$ be a hybrid process and $Y = (X, \alpha)$ its solution. Assume the jump rates do not depend on the continuous state. Let $h > 0$ be a sampling time and $N(x) = E(C)$ the normal distribution with mean $E$ and covariance $C$. Introduce terms

$$\Gamma(t, s) = 0 \int_0^{E(q_i)(t-m)} G(q_i)G(q_i)^\top e^{F(q_i)(t-m)} \ dm,$$

$$\Omega_{\lambda_j}(s) = (\lambda_j - \lambda_j) e^{(\lambda_j - \lambda_j)(t-s)}$$

and for $x \in \mathbb{R}^m$ define $\lambda_j(x) = \frac{\partial}{\partial x_i} \lambda_j(x)$. Then, given $(q, x) \in D$ and a measurable set $(q', A)$, for the resulting DTMP $\tilde{Y} = (\tilde{X}, \tilde{\alpha})$ it holds that
where and $q$ is obtained as the probability of making a jump once, multiplied by the probability associated to paths with more than one jump during $h$. If we assume that the system makes at most one jump during $h$, then, because of the memory-less property of the SDE, during this interval $X$ evolves according to equations as in Lemma 1 and specific to location $q_i$. As a consequence, at time $h$, $X$ is normally distributed, with mean $e^{F(q_i)h}x$ and variance $\Gamma(i,h)$. Instead of the system jumps once within $[0,h]$, after marginalizing over the jump time and the state where this event happens, we end up with a linear Gaussian model [9]. This process is still Gaussian with mean and covariance matrix that can be derived from the equations in Lemma 1. Finally, parameter $\epsilon$ takes into account the probability associated to paths with more than one jump within $[kh, (k+1)h]$; based on the provided upper bound on $\epsilon$, it is clear that the probability of such event becomes negligible as $h$ gets small enough.

The discrete kernel $T^d(q_j, x, q_i)$ has a much simpler derivation. If we assume that the system makes at most one jump during $h$, then the probability that $q_j = q_i$ amounts to the probability that the system does not jump within $[0,h]$. Instead, for the condition $q_j \neq q_i$ the resulting probability is obtained as the probability of making a jump once, multiplied by the probability of jumping to the specific state $q_j$.

From $T^c$ and $T^d$, it is easy to calculate the following transition kernel for $(x, q_i) \in D$ and a measurable set $(A, q_i) \subseteq D$

\[
T^c(A, x, q_i) = \sum_{x \in A} N(x, q_i, h, \Gamma(i,h)) \cdot d\bar{x} \cdot e^{-\lambda_j h} + \frac{\lambda_j}{\lambda_i} \cdot e^{-\lambda_i h} \cdot d\bar{x} + \epsilon \quad \text{if } q_i = q_j,
\]

and

\[
T^d(q_j, x, q_i) = \begin{cases} 
  e^{-\lambda_j h} + \epsilon & \text{if } q_i = q_j, \\
  \lambda_i h - e^{-\lambda_i h} \cdot \frac{\lambda_j}{\lambda_i} + \epsilon & \text{if } q_i \neq q_j,
\end{cases}
\]

where

\[
E_{q_i,x,s}^{\mathcal{H}}(s) = e^{F(q_i)S} e^{F(q_i)(h-s)} x,
\]

\[
C_{q_i,x,s}^{\mathcal{H}} = e^{F(q_i)S} \Gamma(i,s) e^{F(q_i)(h-s)} T + \Gamma(j,h-s),
\]

\[
0 \leq \epsilon \leq 1 - e^{-\lambda_j h} - \lambda_i h - e^{-\lambda_i h}.
\]

In this section we quantify the approximation level introduced by the discretisation procedure. More precisely, we characterize the error associated to the computation of reachability properties with the DTMP over a discrete set of sampling points, with sampling time $h > 0$: by deriving formal error bounds, we show uniform convergence as $h \to 0$.

\textbf{Assumption 2.} Assume that the target set $S$ is independent of the locations, namely select $S \subseteq \mathbb{R}^m$ so that $S = \bigcup_{q \in \mathbb{Q}} q \times S$.

Note that, although this assumption limits the class of properties that can be expressed, there are many applications where target sets independent of discrete locations are of great interest. For example, in the context of controlling room temperature in smart buildings, often the focus is on checking the temperature (continuous variable) regardless of the discrete state of the thermostats (locations). At the end of the section we briefly discuss how this assumption can be relaxed.

Let $H = (\mathbb{Q}, K, F, G, W, \Lambda)$ be a hybrid system and $Y = (X, \alpha)$ its solution. Let $I \subseteq \mathbb{R}_{\geq 0}$ be a finite time interval. For any $q \in \mathbb{Q}$, call $X_q$ the solution of the SDE:

\[
dX_q(t) = F(q)X_q(t)dt + G(q)dB(t).
\]

In this section we assume that $X_q$ is a uni-dimensional, zero mean Gaussian process (GP). In the next subsection, we show how to generalize the results derived here for general GPs and multi-dimensional processes.

$X_q$ is almost surely bounded within the interval $I$ by Assumption 1. Set $h = \min\left\{ \frac{2^n}{2\sqrt{K}r^2} \right\}$ and $\epsilon_n = 2^{-n}$, where $n \in \mathbb{N}$, and $K_d$ is a constant such that for any $t_1, t_2 \in I$

\[
\max_{q \in Q} \left| d_q(t_1, t_2) \right| \leq K_d \cdot |t_2 - t_1|,
\]

where $d_q$ is a pseudometric defined as

\[
d_q(t_1, t_2) = \mathbb{P}\left[ (X_q(t_2) - X_q(t_1))^2 \right],
\]

and $K \geq 12$ is the universal constant in the Dudley’s metric entropy integral [32]. Fix a set of sampling times $\Sigma = \{t_1, ..., t_{|\Sigma|}\}$, with step distance $h$. Call

$S^\ast = \{ x \in S : |x - \delta S| \geq \epsilon_n \}$,

where $\delta S$ is the boundary of $S$ and $| \cdot |$ is the Euclidean distance metrics between a point and a set. Define the events

$A^n = \{ \forall t_n \in \Sigma, X(t_n) \in S^{\ast n} \}$

and

$B = \{ \exists t \in [0, T] \text{ s.t. } X(t) \not\in S \}.$

As $S \subseteq \mathbb{R}$, we have that $P_{safe}(Y, S, I) = P_{safe}(X, S, I)$, where $P_{safe}(X, S, I)$ is the probability that the continuous component $X$ of $Y$, stays in $S$ during $I$. It is easy to see that

$P_{safe}(X, S, I) = \lim_{n \to \infty} P(A^n \land B^n)$.

For a finite $n > 0$, $P(A^n \land B^n)$ is the lower bound for the safety probability computed on $S$, since it requires the system to be inside $S^{\ast n} \subseteq S$ at sampling times in $\Sigma$. Notice that
for $n$ big enough $S$ and $S^n$ become indistinguishable. As a consequence, we can compute the reachability on $S^n$ instead of $S$, assuming $n$ is big enough. Let us define as $P(A^n)$ the reachability probability computed considering only the discrete times in $\Sigma$.

**Theorem 2.** Under Assumption 1, it holds that for all $n \geq 3$ and over a finite time interval $I \subseteq \mathbb{R}_{\geq 0}$

$$P(A^n) \geq P(A^n \land B^n) \geq P(A^n) \cdot \left(1 - \frac{I}{h} \exp(-\frac{2^n - 2^n+1}{2})\right)^n,$$

where $h = \min\{\frac{2^{-n}}{2K+K_d}, 2^{-n}\}$.

**Corollary 1.** Under Assumption 1, it holds that

$$\lim_{n \to \infty} P(A^n \land B^n) = \lim_{n \to \infty} P(A^n).$$

Theorem 2 guarantees that for any $n \geq 3$, we obtain

$$|P(A^n) - P(A^n \land B^n)| \leq \frac{I}{h} \exp(-\frac{2^n - 2^n+1}{2}).$$

This enables choosing, a priori, a sampling interval $h$ that guarantees meeting a chosen error on the precision. The proof of Theorem 2 is given in the Appendix. Here, we explain the main ideas. The proof of Theorem 2 is based on the fact that, for any $q \in \mathcal{Q}$, $X_q$ is a Gaussian process, which is almost surely bounded in $T$. It is possible to show that the supremum of $X_q$ is still distributed as a Gaussian. Then, the use of the entropy Dudley’s integral [21] allows to bound the supremum of $\mathcal{Y}(\hat{X},\alpha)$, assuming $K_d < 12$ and $K_d = \sup_{i \in \mathcal{I}} (K_d,i)$.

Finally, it is important to stress that Assumption 2 can be relaxed by modifying the time discretization error and including a term encompassing the probability that the system jumps more than once during the time interval $[0,h]$. Moreover, as explained next, Theorem 3 can still be used to get lower bounds of cases where the target set depends on the discrete mode. However, the bounds we obtain can be quite conservative if the target sets corresponding to different modes greatly differ.

**Observation 1.** Consider a hybrid system $\mathcal{H}$ with solution $Y = (X, \alpha)$, where $X$ takes values in $\mathbb{R}^m$ and $\alpha$ takes values in a finite set of discrete states $\mathcal{Q}$. Given a measurable set $S = \cup_{q \in \mathcal{Q}} (q \times S_q) \subseteq \mathcal{D}$, we can define $S' = \cup_{q \in \mathcal{Q}} (q \times S_q)$ and $S' = \cup_{q \in \mathcal{Q}} (q \times S_q)$. Then, we have that for a general time interval $I$, $P_{\text{sa}}(Y,S,I) \geq P_{\text{sa}}(Y,S',I)$. That is, if we need to compute probabilistic safety on a set that depends on the discrete modes, then we can always compute a lower bound of this safety considering a target set that is independent of the locations.

**5. State Space Discretization**

In order to complete the procedure leading to a model where we can numerically compute safety or reachability properties, we introduce a numerical scheme inspired by the results of [2, 23]. The numerical scheme is based on a discrete-time Markov chain (DTMC) approximation of the DTMP that results from the time discretization of the original switching diffusion process $\mathcal{H}$. We discuss convergence results and relative error bounds both of this second (state space) approximation step, and of the combined (time- and state approximation) procedure.

Let $S = \cup_{q \in \mathcal{Q}} (q) \times A_q$ be the safe set, where $A_q \subseteq \mathbb{R}^m$. We assume $S$ to be measurable and compact. Given $dx \in \mathbb{R}_{\geq 0}$, we define the grid

$$\mathcal{G}_{dx} = \cup_{q \in \mathcal{Q}} \cup_{i \in \mathcal{I}} (q \times A_i,q),$$

where $A_i,q$ are pairwise disjoint measurable sets, such that for $q \in \mathcal{Q}$ and $i \in \mathcal{I}$, $A_i,q$ are $\alpha_q$-invariant, and $i \neq j$ for $A_i,q \cap A_j,q = \emptyset$, and

$$A_i,q = \{ x, x' \in A_i : |x - x'| \leq dx \}.$$
point \((q, x) \in \{q\} \times A_{i,q}\). The set of representative points \(S_{dx} = \{(q, x_i) \mid i \in \{1, \ldots, m_q\}, q \in Q\}\) makes up the finite state space of the DTMC, a discrete version of the set \(S\). Let us introduce \(\xi : S \rightarrow S_{dx}\), a map that associates to any \((q, x) \in S\) the corresponding representative point. Similarly, the set-valued map \(\Xi : S \rightarrow S_{dx}\) relates any representative point to the concrete \(A_{i,q}\) partition.

We define the discrete state space \(Z_{dx} = S_{dx} \cup \phi\), where \(\phi\) is a discrete state modeling all the states outside \(S\). Note that the compactness of \(S\) guarantees that \(Z\) is finite. The resulting DTMC is completely characterized by its transition kernel \(T_{dx} : Z_{dx} \times Z_{dx} \rightarrow \mathbb{R}_{\geq 0}\), such that for \(z_1 = (x_1, q_1)\), \(z_2 = (x_2, q_2) \in Z_{dx}\), \(T_{dx}(z_1, z_2)\) describes the probability of going in \(z_1\) in the next discrete step, being in \(z_2\) at the current time. \(T_{dx}\) can be easily computed from kernel \(T\) presented in Section 4 as \(T_{dx}(z_1, z_2) = T(z_1, z_2)\), if \(z_1, z_2 \in S_{dx}\), otherwise.

We introduces \(\mathcal{D}\) as a discrete state modeling all the states outside \(S\). Let \(\mathcal{D}\) be a finite set such that \(\exists \mathcal{D} \subset \mathbb{R}_{\geq 0}\).

5.1 Error Bounds for Space Discretization

Let \(\hat{Y}\) be the discrete-time continuous space hybrid process derived through time discretization of \(Y\), solution of the hybrid process \(\mathcal{H}\) with initial condition \((x, q) \in S\). Call \(Y^D\) the approximated DTMC with state space \(S_{dx}\) and initial condition \((x^D, q^D) = \xi((x, q))\). We show that, for \(I \subseteq \mathbb{R}_{\geq 0}\), under Assumption 1, the property \(P_{safe}(Y^D, S_{dx}, I)\) converges uniformly to \(P_{safe}(Y, S, I)\), which also allows us to derive uniform convergence on the original continuous time stochastic process, and to derive error bounds on the global approximations procedure.

**Theorem 4.** [2] Let \(Y\) be the discrete-time continuous space hybrid process with initial condition \((x, q) \in S\), where \(S\) is a measurable set. Call \(Y^D\) the approximated DTMC with state space \(Z_{dx}\), where \(dx > 0\) is the discretization parameter, and initial condition \((x^D, q^D) = \xi((x, q))\). Then, given \(0, N \in \mathbb{N}\), it holds that

\[
|P_{safe}(Y^D, S_{dx}, N) - P_{safe}(\hat{Y}, S, N)| \leq N \cdot K \cdot dx,
\]

where \(K = m_1 h_1 + L h_2\), with \(L\) the Lebesgue measure of the continuous set \(S\), and \(m_1\) cardinality of the discrete set \(Q\).

Notice that, as \(dx \downarrow 0\), the two probabilities collapse.

6. GLOBAL ALGORITHM AND ERRORS

Using the results in Theorem 4, we can derive the uniform convergence between \(P_{safe}(Y^D, S_{dx}, N)\) and \(P_{safe}(Y, S, I)\) for \(h, dx \rightarrow 0\) and \(N\) discretized version of \(I\).

**Theorem 5.** Let \(Y\) be the solution of a switching diffusion process \(\mathcal{H}\) with initial condition \((x, q) \in S\). Call \(Y^D\) the approximated DTMC, with \(h, dx > 0\) time and space discretization parameters, and with initial condition \((x^D, q^D) = \xi((x, q))\). Then, given \(I = [0, t] \subseteq \mathbb{R}_{\geq 0}\), it holds that:

\[
P_{safe}(Y^D, S_{dx}, \frac{I}{h}) - P_{safe}(Y, S, I) \leq \frac{I}{h} \cdot K dx + e^{-\left(2^n - 2^n + 1\right)} 2^{-n} \text{ for } n \geq 3, \text{ with } K \geq 12 \text{ and } K_d \text{ constant introduced in Section 4}.
\]

**Proof.** By triangular inequality we have

\[
P_{safe}(Y^D, S_{dx}, \frac{I}{h}) - P_{safe}(Y, S, I) + |P_{safe}(\hat{Y}, S, I) - P_{safe}(Y, I, I)|.
\]

The proof results from the application of Theorem 4 and Theorem 2.

**Algorithm 1** Probabilistic safety computation by finite DTMC abstraction

**Require:** \(Y = (X, \alpha)\) solution of \(\mathcal{H}\) with initial condition \((x, q)\), safe set \(S\), finite time interval \(\delta = [0, t]\), and parameters \(dx, h = \min \frac{2^{-n}}{2\sqrt{2\pi K^2} \cdot dx} 2^{-n} ;
1: Select the partition \(\mathcal{D}_{dx} = \bigcup_{q \in Q} \{q\} \times A_{i,q};\)
2: Select the set of representative points, leading to \(S_{dx};\)
3: Define the DTMC \(Y^D\) with state space \(Z_{dx} = S_{dx} \cup \phi\), initial condition \(z_0\) equals to 1 for the entry corresponding to \(\xi((x, q))\) and 0 otherwise, and transition matrix \(P_{dx}\) such that \(P_{dx}(i, j) = T_{dx}(z_i, z_j);\)
4: Compute \(z' = z_0 \cdot P_{dx}\)
5: Return \(P_{safe}(Y, S, I) = 1 - z'(\phi)\) with the error given as \(\frac{I}{h} \cdot (K dx + e^{-\left(2^n - 2^n + 1\right)} 2^{-n})).

In Algorithm 1 we present a numerical routine to compute safety properties over continuous-time hybrid systems. The inputs of the algorithm are \(Y = (X, \alpha)\), solution of the continuous time hybrid process \(\mathcal{H}\) with a given initial condition, a finite time interval \(I\), the sampling time \(h\), the grid parameter \(dx\) and the target set \(S\). (In the case study presented in the next section we consider parameters \(h = 0.1, dx = 0.2\) and \(I = [0, 2]\).) Theorem 5 allows us to compute a bound on the error as a function of parameters \(dx\) and \(h\). Moreover, such parameters can be selected to meet a required precision error. That is, given the maximum error that is tolerated, Theorem 5 returns possible \(h\) and \(dx\) that guarantee such an error. In Lines 1, 2, 3 the algorithm computes the DTMC abstraction from \(Y\) and \(S\), as described in the previous section: \(P_{dx}\) is the transition probability matrix of the resulting DTMC [31], namely \(P_{dx}(i, j)\) describes the probability of going from the discrete state \(z_i\) to the discrete state \(z_j\) at the
next time step. Line 4 computes the transient evolution of the DTMC $Y^D$. This is done by multiplying the initial state $z_0$ for $Pdz \left[ \frac{1}{h} \right]$ times, where $[\frac{1}{h}]$ are the number of discrete steps: $P_{safe}(Y,S,I)$ is just the probability of not being in the sink state $\phi$. A bound on the error is computed using Theorem 5.

7. CASE STUDY

We consider a continuous-time switching diffusion process studied in [1]. The discrete state space is composed of two locations $Q = \{on, off\}$, and the continuous process takes values in $\mathbb{R}^2$, so that the hybrid state space is $D = Q \times \mathbb{R}^2$. The drift is given by the following two matrices $F(on) = -0.6 \ 0.3 \ -0.6 \ 0.15 \ , \ F(off) = -0.35 \ 0 \ 0.1 \ -0.25$.

The continuous dynamics are further affected by a 1-dimensional Wiener process scaled by matrices

$$G(on) = \begin{bmatrix} 0.2 \\ 0.2 \end{bmatrix}, \ G(off) = \begin{bmatrix} 0.3 \\ 0.3 \end{bmatrix}. $$

The Poisson measures are independent of the continuous component of the process and are given by the following rates: $\lambda_{on, off} = 0.41$ and $\lambda_{off, on} = 0.38$. We consider the Borel sigma algebra over $D$ and a measurable set $A$. As the rates are independent of the continuous components, for $A \subseteq \mathbb{R}^2, q_i, q_j \in \{on, off\}$ with $q_i \neq q_j, x \in \mathbb{R}^2$ and $h \in \mathbb{R}_{>0}$ small enough, we have the following transition kernels (see Theorem 1):

$$T^c(A, x, q_i, k) = \begin{cases} N \ {\bar{x}}e^{F(q_i)h} x, \Gamma(i, h) \ dx \cdot e^{-\lambda_i,h} + \\ Z \ Z_h \ \int A \ {\bar{x}}e^{F(q_i)h} x, C_{q_i,x,s}^{H} \cdot \Omega_{i,j,h}(s)ds, \end{cases}$$

$$(\lambda_i, h \ e^{-\lambda_i,h})dx, \ \text{where} \ e^{F(q_i)h} x = e^{F(q_i - s)}e^{F(q_i)(h-s)}x, \ \text{and} \ C_{q_i,x,s}^{H} = e^{F(q_i) - \Gamma(i,s)}e^{F(q_i)(h-s)} + \Gamma(j,h-s),$$

$$T^d(q_j, x, q_i, k) = \begin{cases} e^{-\lambda_i,h} & \text{if } q_i = q_j, \\ \lambda_i,h \cdot e^{-\lambda_i,h} & \text{if } q_i \neq q_j. \end{cases}$$

In order to choose $h$, we need to compute constants $K_1, h_2$, and $h_2$. As the rate coefficients are independent of the continuous components we have $h_2 = 0$. It can be further derived that

$$h_2 \leq \max_{x \in \mathbb{R}^2} \left| \frac{\partial \mathbb{E}^c(x'|x, q_i)}{\partial x} \right|, $$

where $t'$ is the density function of the kernel $T^C$. Further, $K_d$ can be computed as

$$K_d = \max_{q_i \in \{on, off\}, j \in \{1, 2\}} \mathbb{P}(i, h)(j, j),$$

where $\Gamma(i, h)(j, j)$ is the component $(j, j)$ of matrix $\Gamma(i, h)$. Note that $K_d$ is also independent of the continuous component of the process.

In order to demonstrate the soundness of our method we implement Algorithm 1 in Matlab, and compare the numerical implementation with empirical results obtained by simulations. We consider the following safe region

$$\mathcal{S} = x \in \mathbb{R}^2 \text{ s.t. for } i \in \{1, 2\}, -0.2 \leq x_i \leq 1,$$

where $x_i$ is the $i$-th component of vector $x$. We select the time interval $I = [0, 2]$. Firstly, we consider $h = 0.1$ and $dz = 0.2$. For such values, the resulting abstract DTMC is made up of 5184 states. We consider different initial conditions and we compare the safety computed on the abstraction with the same property computed on the original continuous-time model using 1000 simulations. As expected, for any initial condition, the abstraction provides a safety value that is a lower bound of the empirical one. We observe a maximal error of 0.11. Note that, for $h = 0.1$ and $dz = 0.2$, Theorem 5 guarantees a theoretical time discretization error bound of 0.2, and an uninformative space discretization error bound (e.g. > 1). This is because, for small values of $h$, the value of the constant $h_2$ tends to increase, requiring a finer space grid.

In order to increase the precision it is possible to decrease $h$ and $dz$ at the price of more computational effort (a larger DTMC abstraction). To guarantee a theoretical error $\leq 0.1$, we can select $h = 0.03$, which results in a theoretical time discretization error $\leq 4 \times 10^{-4}$. However, for such small $h$, $t'$ has very small variance, rendering $h_2$ large. As a consequence, in order to keep the error small, we would need $dz < 10^{-3}$ and the resulting DTMC would be composed of $> 10^8$ states. This dimensionality issue arises also because we are considering a uniform grid ($dx$ constant). As a consequence, we use the same space resolution both for states with no probability mass and for states with large probability mass, which are the great minority for $h$ small. In fact, as described in the next section, The use of more advanced, adaptive grid techniques [23] would allow us to meet the given precision with a much smaller resulting DTMC – this is targeted as future work.

8. CONCLUSIONS

We have presented a novel and formal approach to compute probabilistic reachability (and dually safety) for continuous time hybrid processes with no guards and no resets, and with continuous dynamics that can be described by linear stochastic differential equations. We have considered an approach based on space and time discretization of the original process, and derived uniform convergence of the algorithm, as well as error bounds that can be used to tune and control the approximation error.

The main contributions of the paper are the characteriza-

the time discretization of such pro-

the error bound for the time discretization pro-

Finding formal bounds for the time discretization of stochastic hybrid processes has been an open problem, and, to our knowledge, only limited to results of weak conver-

the bound for uni-dimensional target sets, and then have shown how to extend it to multidimensional processes.

For the space discretization we have considered an ap-

the work in [2]. Although formally correct, this approach in combination with time discretization may result in large DTMC abstractions. In fact, as shown in a case study, the diameter of each grid location tends to grow as the sam-

the sampling time of the time discretization process decreases. A much better solution would be to consider adaptive grid-

these would be extremely beneficial, since, when the sampling time is small, the distribution of
9. PROOFS

Proof of Theorem 1. We show the derivation of the continuous kernel. Note that the discrete kernel can be derived similarly. By definition we have

\[
T^c(A, x, q_i) = \int_A t^c(\bar{x}|x, q_i)d\bar{x},
\]

where \(t^c(\bar{x}|x, q_i)\) is the density function of \(X\), continuous component of \(Y\), assuming \(X(0) = x\) and \(\alpha(0) = \lambda_i\). We define \(\text{Num}_i^h = k\) as the event such that \(\alpha\), discrete component of \(Y\), jumps \(k\) times between \([0, h]\). By marginalizing with respect to the number of times that \(\alpha\) jumps during \([0, h]\), we have

\[
\begin{align*}
\int_A t^c(\bar{x}|x, q_i) d\bar{x} &= \text{Z} \\
&= \int_A \left( \text{Prob}(\text{Num}_i^h = 0) \cdot \text{Prob}(\text{Num}_i^h = 1) \right) \cdot \text{Prob}(\text{Num}_i^h = 1 | x, q_i) + \epsilon d\bar{x} \\
&= 1 - \text{Prob}(\text{Num}_i^h = 1 | x, q_i) \\
&= e^{-\lambda_i h}.
\end{align*}
\]

The first term is a linear Gaussian model. This class of models has been extensively studied in literature [9]. More specifically, it has a Gaussian distribution whose variance and expectation can be derived from Lemma 3. As a consequence, we have

\[
\begin{align*}
t^c(\bar{x}|x, q_i, \text{Jump}_{i,j}) &= \mathcal{N}(\bar{x}|E, C), \\
E &= e^{F(q_i)} + e^{F(q_i)(h-s)} x, \\
C &= e^{F(q_i)^h} (\Gamma(i, s) - e^{F(q_i)} + \Gamma(j, h - s)),
\end{align*}
\]

Proof of Theorem 2. For each \(q \in Q\), \(X_q\) is an almost surely bounded and uniformly continuous GP in \(I\), time interval of interest. This guarantees the existence of a sequence \(\{\delta_n\}\) with \(\delta_n \to 0\) such that \(\phi(\delta_n) \leq 2^{-n}\) (see Theorem 2.1.3 of [5]), where

\[
\phi(\delta_n) = E \max_{q \in Q} \sup_{s, s' \in I, |x - x'| \leq \delta_n} (X(s) - X(s')).
\]

Set \(h = \min\{\delta_n, 2^{-n}\}\) and \(\epsilon_n = 2^{-\frac{n}{2}}\). For a set of a sampling times \(\Sigma = \{t_1, \ldots, t_{\beta(n)}\}\), with step distance \(h > 0\), call \(S \subseteq \mathbb{R}^n\) the safe set and,

\[
\begin{align*}
S^n &= \{x \in S : |x - \partial S| \leq \epsilon_n\}, \\
A^n &= \{\forall t_i \in \Sigma, X(t_i) \in S^n\}
\end{align*}
\]

and

\[
B = \{\exists \ i \in I \text{ s.t. } X(t) \notin S\}.
\]

Using the rules of probability we get

\[
P(A^n \land B) = P(A^n) \cdot (1 - P(B|A^n))
\]

By definition of probability we have that

\[
\begin{align*}
0 &\leq P(B|A^n) \\
&\leq P(\exists t_i \in \Sigma, s.t. \sup_{t \in [s_i, s_i+h]} X(t) - X(i) > \epsilon_n) \\
&\leq P(\exists t_i \in \Sigma, s.t. \sup_{t \in [s_i, s_i+h]} X(t) - X(i) > \epsilon_n) \\
&\leq \frac{\epsilon_n}{\sigma} P \sup_{t \in [s_i, s_i+h]} (X(t) - X(i)) > \epsilon_n,
\end{align*}
\]

with \(X(t_i) \in S^n\).

In order to bound \(P(\sup_{t \in [s_i, s_i+h]} (X(t) - X(i)) > \epsilon_n)\) we need to take into account that during \([s_i, s_i+h]\) the discrete state may hit a transition. However, there is no reset for the continuous components. As a consequence, it is enough to assume that, during \([s_i, s_i+h]\), \(X\) always evolves according to the "worst" behaving \(X_q\). Then, being \(X_q\) a GP, we can make use of the Borell's bound [4, 5]. Given an interval \(I\) and a centered and bounded Gaussian process \(X_q\) with \(\sigma_1 = \sup_{t \in I} \sigma(t)\), supremum of the standard deviation of the process, the Borell bound guarantees that

\[
\sup_{t \in I} X_q(t) > u \leq \exp\left(-\frac{u - E[X_q(t)]^2}{\sigma(t)^2}\right).
\]

Applying this result to our case for intervals of the type \([s_i, s_i+h]\), and for \(u = \epsilon_n = 2^{-\frac{n}{2}}\), where \(n \geq 3\) we have that
\[ P \sup_{t \in [t_i, t_i + h]} (X(t) - X(t_i)) \leq \epsilon_n \]

\[ \left| \sum_n \right| \exp \left( \frac{z - 2 - z^2}{2 - z^n} \right) \leq \left| \sum_n \right| \exp \left( \frac{z^2 - 2 + z}{2 - z^n} \right) \]

At this point, the last, and non-trivial, step in order to derive our convergence results and relative error bounds is to show that

\[ \lim_{n \to \infty} \left| \sum_n \right| \exp \left( \frac{z^2 - 2 + z}{2 - z^n} \right) = 0. \]

In fact, as

\[ 0 \leq P(B|A^n) \leq \left| \sum_n \right| \exp \left( \frac{z^2 - 2 + z}{2 - z^n} \right), \]

this would guarantee that

\[ P_{\text{safe}}(X, S, I) = \lim_{n \to \infty} P(A^n \land B^n) = \lim_{n \to \infty} P(A^n). \]

To do that, it is sufficient to show that \( h = \frac{z^2 - 2}{z^n} \), for some constant \( C \). In fact, this implies \( \left| \sum_n \right| = \frac{C}{z^n} \).

Recall that we chose \( h \) such that for all \( t_i \in \Sigma_n \),

\[ E \sup_{t \in [t_i, t_i + h]} (X(t) - X(t_i)) \leq 2^{-n}. \]

As a consequence, it is enough to take \( h \) as the greatest interval smaller than \( 2^{-n} \) such that this condition is verified. For \( t_i \in \Sigma_n \), call \( X_t = X(t) - X(t_i) \). We can now make use of the Dudley integral (or entropy integral) [4], which guarantees that for \( t_i \in \Sigma_n \),

\[ E \left[ \sup_{t \in [t_i, t_i + h]} \left| \sum_n \right| \exp \left( \frac{z^2 - 2 + z}{2 - z^n} \right) \right] \leq K \frac{\ln(N([t_i, t_i + h], d, \epsilon))}{d \epsilon}, \]

where \( K \geq 12 \) is a constant and \( d \) is a pseudo-metric defined as

\[ d(t, t + dt) = P(E[(X(t) + dt) - X(t)]^2], \]

\[ N([t_i, t_i + h], d, \epsilon) \]

represents the smallest number of balls of radius \( \epsilon \), which covers \([t_i, t_i + h]\), under metric \( d \), where \( \text{diam}([t_i, t_i + h]) \) is defined as

\[ \text{diam}([t_i, t_i + h]) = \sup_{s', s \in [t_i, t_i + h]} d(s', s) \]

and with our assumptions, it is possible to show that there exists a constant \( K_d \) such that

\[ d(t, t + h) \leq K_d \cdot h \]

Moreover, for \( \tilde{T}_i = [t_i, t_i + h] \subseteq \mathbb{R}_{\geq 0} \) we have

\[ N(\tilde{T}_i, d, \epsilon) \leq \frac{K_d h}{2\epsilon} + 1, \]

This can be easily understood thinking at the geometry of the problem. As a consequence, we have

\[ E \sup_{t \in \tilde{T}_i} (X(t)) \leq K \frac{Z \sqrt{2 - n} \int_0^{2^{2-n-1}} S \ln \left( \frac{K_d h}{2\epsilon} + 1 \right) d\epsilon}{2}. \]

Now, our property is satisfied if we chose \( h \) such that

\[ K \frac{Z \sqrt{2 - n} \int_0^{2^{2-n-1}} S \ln \left( \frac{K_d h}{2\epsilon} + 1 \right) d\epsilon}{2} \leq 2^{-n}. \]

The integral inequality we need to solve cannot be solved analytically. However, as \( K_d h > 0 \), we can write

\[ \int_0^{2^{2-n-1}} S \ln \left( \frac{K_d h}{2\epsilon} + 1 \right) d\epsilon \leq K \frac{K_d}{2\epsilon} \frac{Z \sqrt{2 - n}}{2} \frac{r}{\epsilon} \frac{d\epsilon}{2^{2-n}}. \]

Asking for this quantity to be smaller than \( 2^{-n} \), we obtain the following bound for the sampling time \( h \):

\[ h \leq \min \left( \frac{2^{-n}}{2\sqrt{2K^2K_d}}, 2^{-n} \right). \]

\[ \square \]

**Acknowledgments**

This work has been supported by the Royal Society Professorship (L. Cardelli), the Czech Grant Agency grants GA16-24707Y (M. Češka), the EU-FET project QUANTICOL (nr. 600708) and by FRA-UniTS (L. Bortolussi).

**10. REFERENCES**


